

On computational complexity of Khovanov homology

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Abstract

Computing the Jones polynomial of general link diagrams is known to be $\#P$ -hard, while restricting the computation to braid closures on fixed number of strands allows for a polynomial time algorithm. We investigate polynomial time algorithms for Khovanov homology of braids and show that for 3-braids there is one. In contrast, we show that Bar-Natan's scanning algorithm runs in exponential time when restricted to simple classes of 3-braids. For more general braids, we obtain that a variation of the scanning algorithm computes the Khovanov homology for a bounded set of homological degrees in polynomial time. We also prove upper and lower bounds on the ranks of Khovanov homology groups.

1 Introduction

The original construction of Khovanov presented a link invariant, Khovanov homology, which is theoretically computable for any link diagram [Kho00], but the number of steps required to do so grows exponentially in the number of crossings. A few years later, Bar-Natan's work on tangles made it practically possible to compute Khovanov homology for links with 50 or more crossings [BN07], which would be impossible from the definition. This advancement in algorithms has opened up new avenues in 4-dimensional topology: In [FGMW10] a possible (yet to be realized) attack for disproving the smooth 4-dimensional Poincaré conjecture was formulated and in [Pic20] a computer calculation helped to prove that the Conway knot is not smoothly slice.

The Jones polynomial $J_L(q)$ of a link L can easily be recovered from its Khovanov homology $\text{Kh}(L)$ and hence Khovanov homology is at least as hard to compute as the Jones polynomial. By reducing the computation of the Jones polynomial of a link diagram L to the evaluation of the Tutte polynomial of a related planar graph, Jaeger, Vertigan and Welsh showed that $J_L(q)$ is $\#P$ -hard to compute with respect to the number of crossings on L [JVW90]. It follows that the existence of a polynomial time algorithm which, given a link diagram L , computes the Khovanov homology $\text{Kh}(L)$ is very unlikely. On the other hand, Morton and Short showed that the Jones polynomial of a closed braid on a fixed number of strands can be computed in polynomial time [MS90]. This led to the following conjecture which provides the main motivation for this article.

Conjecture 1.1 (Przytycki, Silvero [PS24]). *Computing Khovanov homology of a closed braid with fixed number of strands has polynomial time complexity with respect to the number of crossings.*

The state-of-the-art for computing Khovanov homology is given by Bar-Natan's scanning algorithm [BN07]. However, there exist relatively simple links which make this algorithm very slow.

Theorem 1.2. *The scanning algorithm and the divide-and-conquer algorithm of [BN07] run in exponential time with respect to the number of crossings in a link diagram, even when restricted to positive 3-braids, or to alternating 3-braids.*

The reason for the slow runtime of these algorithms is that there are such 3-braids whose total Betti numbers grow exponentially with respect to the number of crossings. Any algorithm, which works with a basis on the chain complex, will therefore necessarily have an exponential runtime. As with classical algorithms, the large Betti numbers also played a significant role in the recent study of quantum algorithms for Khovanov homology [SRZ⁺25]. The exponential Betti numbers make it impossible for the main algorithm of [SRZ⁺25] to accurately estimate the Khovanov homology of 3-braids in polynomial time. For other standard quantum homology algorithms [SRZ⁺25, §3] large Betti

numbers seem to be an asset although it remains unclear how apply them to Khovanov homology. Despite the exponential Betti numbers, we give an affirmative answer to Conjecture 1.1 in the case of braids on 3 strands.

Theorem 1.3. *Conjecture 1.1 is true for 3-braids.*

The algorithm given in Section 3 to prove Theorem 1.3 is based on the earlier work of the second author [Sch25b], which showed that the integer Khovanov homology of a closed 3-braid either is described by a formula in [CLSS22], or decomposes roughly as a direct sum of Khovanov homologies of a torus link $T(3, 3k)$ and an alternating link. The homology groups $\text{Kh}(T(3, 3k))$ have also been computed in [CLSS22] while the Khovanov homology of the alternating links can be obtained from their Jones polynomials and the signatures. These alternating links are themselves braid closures, so the polynomial time algorithm of Morton and Short can be used.

The Khovanov complex $\text{CKh}(L)$ of a link diagram L is a bigraded complex of finitely many finite rank \mathbb{Z} -modules. Denote the minimal square of \mathbb{Z}^2 gradings, which supports $\text{CKh}(L)$ by $[i_{\min}, i_{\max}] \times [q_{\min}, q_{\max}]$. As evidence for Conjecture 1.1 Przytycki and Silvero gave a polynomial time algorithm, which computed $\text{Kh}^{*, q_{\min}}(L)$ (and by symmetry $\text{Kh}^{*, q_{\max}}(L)$) for 4-braids. The construction of $\text{CKh}(L)$ gives $i_{\min} = -n_-(L)$ and $i_{\max} = n_+(L)$ where $n_+(L)$ and $n_-(L)$ denote the number of positive and negative crossings of L . We show that the Khovanov homology at any number of homological gradings near i_{\min} and i_{\max} for braids on any number of strands can be obtained in polynomial time.

Theorem 1.4. *For every $k, t \geq 0$ there is an algorithm $\mathcal{B}_{k,t}$ which takes in a braid b on t strands and outputs the integral Khovanov homology $\text{Kh}^{i,*}(L_b)$ of the braid closure L_b in homological degrees $i \leq -n_-(L_b) + k$ and $i \geq n_+(L_b) - k$. The algorithms $\mathcal{B}_{k,t}$ run in polynomial time with respect to the length of b .*

The algorithms $\mathcal{B}_{k,t}$ are only a minor variation of Bar-Natan's scanning algorithm. While scanning, the intermediate complexes of tangles can be truncated without loosing information relevant to the extremal homological degrees. A similar trick was used in a fast algorithm for computing s -invariants [Sch21]. Even in the non-truncated degrees, the Khovanov hypercube complex may have exponential rank so we have to ensure that the Gaussian elimination of the scanning algorithm cancels out 'most' of the summands. Since the intermediate complexes contain integer coefficients c , which need to be stored in the memory, we also need some polynomial p for which $\log_2(c) \leq p(n(L_b))$ for all c we process. As a by-product of our proof we obtain the following upper bound where $n(L) = n_+(L) + n_-(L)$:

Proposition 1.5. *For a connected link diagram L and a field \mathbb{F}*

$$\dim_{\mathbb{F}}(\text{Kh}^{i,*}(L; \mathbb{F})) \leq 2 \binom{n(L)}{i + n_-(L)}.$$

The bounds of Proposition 1.5 are asymptotically strict in the following sense:

Proposition 1.6. *Let L_t be the braid closure of $(\sigma_1 \sigma_3 \sigma_2^4)^t \sigma_1 \sigma_3$. For every $k \geq 0$ there exists $\varepsilon > 0$ for which*

$$\liminf_{t \rightarrow \infty} \left(\frac{\text{rank}(\text{Kh}^{n(L_t)-k,*}(L_t))}{\binom{n(L_t)}{n(L_t)-k}} \right) \geq \varepsilon.$$

Proposition 1.6 is a non-vanishing result so proving it requires us to generate a lot of distinct homology cycles. To do so, we take the Khovanov complex of L_t and employ a Morse matching on it from the earlier work of the first author [Kel25]. This creates a chain homotopy equivalent Morse complex, where sufficiently many copies of \mathbb{Z} split off. The same asymptotic lower bound is also proven for odd Khovanov homology.

Is Conjecture 1.1 true? As stated in [PS24], establishing Conjecture 1.1 would be a game changer in stating and testing conjectures about Khovanov homology. While Theorem 1.3 gives hope for this conjecture to be true, its proof very much uses the particular structure of 3-braids and does not seem to generalize easily to more strands. Furthermore, it reveals that the Khovanov homology of any closed 3-braid roughly decomposes into a part coming from an alternating braid word and a part coming from a torus link word. Already in the case of 4-braids the general structure seems much more complicated. For example, [MPS⁺18] produced positive 4-braids containing torsion of order 3

in their Khovanov homology, as well as others that contain torsion of order 2^k for k up to 23. Also, computations show that 4-strand torus links have a much more complicated Khovanov homology than their 3-stranded counterparts.

Disproving Conjecture 1.1 would amount to constructing a reduction to a known NP-hard problem. The authors are not aware of any promising candidates for such reduction.

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2 Preliminaries

In this section we fix the notation for the Khovanov complex and a few variations, as well as our conventions for braids and algebraic discrete Morse theory.

By a *Frobenius algebra* A over a commutative ring R we mean a commutative ring A that contains R as a subring that comes with an R -module map $\varepsilon: A \rightarrow R$ and an A -bimodule map $\Delta: A \rightarrow A \otimes_R A$ that is co-associative and co-commutative, and such that $(\varepsilon \otimes \text{id}) \circ A = \text{id}$.

In the case that A is free of rank 2 over R , Khovanov [Kho06] showed that for a link diagram L one can define a cochain complex $C(L)$ over R whose homology is a link invariant. The Frobenius algebra leading to Khovanov's original link homology [Kho00] is given by $A = \mathbb{Z}[X]/(X^2)$, with $\varepsilon(1) = 0$, $\varepsilon(X) = 1$, and $\Delta(1) = 1 \otimes X + X \otimes 1$. Here A can be given a q -grading by $|1|_q = 1$ and $|X|_q = -1$, which turns the *Khovanov complex* $\text{CKh}(L)$ into a bigraded cochain complex. We write $\text{CKh}^{i,j}(L)$ to indicate the homological grading and the q -grading, and $\text{Kh}^{i,j}(L)$ for the corresponding homology groups. We may also omit the gradings, or write $\text{Kh}^{i,j}(L; R)$ if we want to indicate coefficients in a ring R .

For tangles T there is also a Khovanov complex, we follow mainly Bar-Natan's dotted version [BN05]. Let $B \subset \mathbb{R}^2$ be a compact surface and $\dot{B} \subset \partial B$ an oriented 0-dimensional compact manifold bordant to the empty set. We denote by $\text{Cob}_\bullet(B, \dot{B})$ the category whose objects are compact smooth 1-dimensional submanifolds $S \subset B$ intersecting ∂B transversely in $\partial S = \dot{B}$. Morphisms between objects S_0 and S_1 are isotopy classes of dotted cobordisms between $S_0 \times \{0\}$ and $S_1 \times \{1\}$ embedded in $B \times [0, 1]$ and which are a product cobordism near $\dot{B} \times [0, 1]$. Here dotted means that there are finitely many specified points in the interior of the cobordism, which can move freely there.

If R is a commutative ring, we denote by $\text{Cob}_{\bullet/l}^R(B, \dot{B})$ the additive category whose objects are finitely generated free and based R -modules where basis elements are objects $q^j S$ with S an object of $\text{Cob}_\bullet(B, \dot{B})$ and $j \in \mathbb{Z}$. Here q^j stands for a shift in the q -grading. Morphisms are given by matrices (M_{nm}) with each matrix entry M_{nm} an element of the free R -module generated by the morphism set between S_n and S_m in $\text{Cob}_\bullet(B, \dot{B})$, modulo the relations in Figure 1. The category $\text{Cob}_{\bullet/l}^R(B, \dot{B})$ admits a delooping isomorphism, see Figure 2, which is at the heart of many of our computations.

Figure 1: The dotted relations of $\text{Cob}_{\bullet/l}^Z(B, \dot{B})$. We refer to the two leftmost as *sphere relations*, the middle one as *double dot relation* and the right one as *neck cutting relation*.

Given an oriented crossing $c \in \{\times, \text{X}\}$ we denote by $\llbracket c \rrbracket$ the tangle complex of c , or Bar-Natan complex of c , defined as

$$\begin{aligned} \llbracket \times \rrbracket &= \dots \longrightarrow 0 \longrightarrow u^0 q^1) \xrightarrow{\text{saddle}} u^1 q^2 \text{X} \longrightarrow 0 \longrightarrow \dots \\ \llbracket \text{X} \rrbracket &= \dots \longrightarrow 0 \longrightarrow u^{-1} q^{-2} \text{X} \xrightarrow{\text{saddle}} u^0 q^{-2}) \longrightarrow 0 \longrightarrow \dots \end{aligned}$$

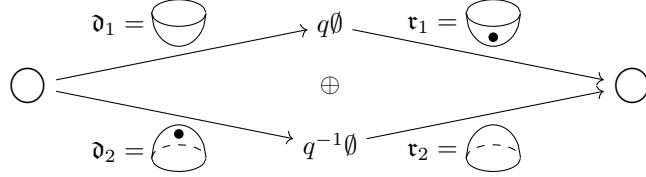


Figure 2: The local delooping isomorphism $\begin{bmatrix} \mathfrak{d}_1 \\ \mathfrak{d}_2 \end{bmatrix}$ its inverse $\begin{bmatrix} \mathfrak{r}_1, \mathfrak{r}_2 \end{bmatrix}$.

over the additive category $\mathcal{C}ob_{\bullet/l}^R(B, \dot{B})$, where \dot{B} are four points with orientation induced by c . Here u^i denotes the homological degree and q^j the q -degree.

The categories $\mathcal{C}ob_{\bullet/l}^R(B, \dot{B})$ compose well, see [BN05, §5], allowing us to form tensor products between complexes $\llbracket c_1 \rrbracket$ and $\llbracket c_2 \rrbracket$, which in turn leads to tangle complexes $\llbracket T \rrbracket$ over $\mathcal{C}ob_{\bullet/l}^R(B, \dot{B})$ for any tangle diagram T embedded in B with endpoints given by \dot{B} .

Given a word w in an alphabet $\{\sigma_i, \sigma_i^{-1} \mid 1 \leq i \leq n-1\}$ for some $n \geq 2$, we get a tangle T_w by using a standard tangle for each σ_i^\pm as in Figure 3 and stacking them on top of each other. This gives rise to a braid on n strands which can be closed to a link.

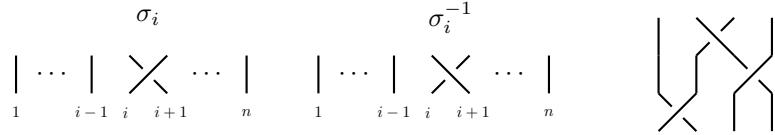


Figure 3: The tangles for σ_i , σ_i^{-1} , and $\sigma_1\sigma_3\sigma_2^{-1}$.

We note that our conventions for braid diagrams lead to the mirror of the braid diagrams in [Mur74]. The reason we use our convention is that we want positive braid words to give rise to links whose Khovanov homology is concentrated in non-negative homological degrees.

Also, the tangle diagram T_w fits in a rectangle with n points on the bottom and n points on the top. We simply write $\mathcal{C}ob_{\bullet/l}^R(B_n^n)$ for the corresponding category over which the Khovanov complex is defined. In later sections, the compact surface $B \subset \mathbb{R}^2$ is simply a disc with \dot{B} consisting of $2n$ points. We will then write $\mathcal{C}ob_{\bullet/l}^R(2n)$ for $\mathcal{C}ob_{\bullet/l}^R(B, \dot{B})$.

2.1 Gaussian elimination and algebraic discrete Morse theory

The following standard lemma is a key tool for simplifying combinatorial chain complexes, both from a theoretical and an algorithmic viewpoint.

Lemma 2.1 (Gaussian elimination). *Let (C, d) be a chain complex over an additive category \mathfrak{C} . Assume that for n there are decomposition of the chain spaces $C^n = C_1^n \oplus C_2^n$ and $C^{n+1} = C_1^{n+1} \oplus C_2^{n+1}$ and differentials d^{n-1} , d^n , and d^{n+1} so that*

$$C = \dots \rightarrow C^{n-1} \xrightarrow{\begin{bmatrix} \alpha \\ \beta \end{bmatrix}} C_1^n \oplus C_2^n \xrightarrow{\begin{bmatrix} \varphi & \delta \\ \gamma & \varepsilon \end{bmatrix}} C_1^{n+1} \oplus C_2^{n+1} \xrightarrow{\begin{bmatrix} \mu & \nu \end{bmatrix}} C^{n+2} \rightarrow \dots$$

If $\varphi: C_1^n \rightarrow C_1^{n+1}$ is an isomorphism, then C is chain homotopy equivalent to the complex C' which in homological degrees $[n-1, n+2]$ is defined by

$$C' = \dots \rightarrow C^{n-1} \xrightarrow{\beta} C_2^n \xrightarrow{\varepsilon - \gamma \varphi^{-1} \delta} C_2^{n+1} \xrightarrow{\nu} C^{n+2} \rightarrow \dots$$

and outside these degrees C' agrees with C .

Applying Gaussian elimination once on a large chain complex C often does not cut it and one needs to iterate Lemma 2.1 in order to obtain a reasonable description of C . Since the Gaussian elimination

introduces a new morphism between the remaining summands, it is not immediately obvious whether an isomorphism of C continues to be a cancellable isomorphism after some number of other Gaussian eliminations have been performed. Algebraic discrete Morse theory [Skö06] is a framework which provides a sufficient condition for iteratively applying Lemma 2.1 and a description of the resulting complex.

A *based chain complex* is a chain complex (C, d) over an additive category \mathfrak{C} and a fixed direct sum decomposition $C^i = \bigoplus_k C_k^i$ on every chain space C^i . The decomposition of the chain spaces also gives out a decomposition of the differentials: we call the maps

$$d_{C_k^{i+1}, C_j^i} : C_j^i \rightarrow C_k^{i+1}, \quad d_{C_k^{i+1}, C_j^i} = \pi d^i \iota$$

matrix elements of C , where ι and π are the canonical inclusions and projections associated to C_j^i and C_k^{i+1} . The based complex C induces a directed graph $G(C) = (V, E)$, whose vertices V are the direct summands C_k^i and whose directed edges E are the non-zero matrix elements of C . Reversing a subset of edges $M \subset E$ gives out another graph $G(C, M) = (V', E')$ with $V' = V$ and

$$E' = (E \setminus M) \cup \{b \rightarrow a \mid (a \rightarrow b) \in M\}.$$

We call M a *Morse matching* on C , if

1. M is finite.
2. M is a matching, i.e., its edges are pairwise non-adjacent.
3. For every edge $f \in M$, the corresponding matrix element f is an isomorphism in \mathfrak{C} .
4. $G(C, M)$ has no directed cycles.

The graph $G(C, M)$ can be viewed as a category, whose objects and morphisms are the vertices and directed paths of $G(C, M)$. This allows us to define a functor $R: G(C, M) \rightarrow \mathfrak{C}$ on vertices by $R(C_k^i) = C_k^i$. On single edges, we set $R(f) = f$ if $f \notin M$ and $R(g) = -g^{-1}$, if $g \in M$ and on longer paths we extend R functorially. If M is a Morse matching on C , we can define the based *Morse complex* $(M(C), \partial)$, whose direct summands are the summands of C which are not matched by M . The differentials ∂^i are described by their matrix elements with

$$\partial_{B,A}^i = \sum_{p \in \{\text{paths: } A \rightarrow B\}} R(p)$$

where A and B are unmatched cell of C^i and C^{i+1} respectively.

Theorem 2.2 (Sköldberg [Skö06]). *If M is a Morse matching on C , then $C \simeq M(C)$.*

Sketch of proof. Inductively use Lemma 2.1 on the isomorphisms of M . Condition 4 ensures that the $k+1$:th isomorphism remains an isomorphisms after k cancellations. \square

From the four conditions of Morse matching (2) is often easy to prove and it greatly simplifies the combinatorics of $G(C, M)$. If (2) holds, then no path can descend two homological degrees and to prove (4) it suffices show that there exist no loops which zig-zag between two adjacent homological degrees. Similarly, all paths that contribute to the differentials ∂ of the Morse complex $M(C)$ have to zig-zag between adjacent homological degrees.

3 Khovanov homology of 3-braids in polynomial time

The braid group on 3 strands is given by $\mathbf{B}_3 = \langle \sigma_1, \sigma_2 \mid \sigma_1 \sigma_2 \sigma_1 = \sigma_2 \sigma_1 \sigma_2 \rangle$. We denote the element $\Delta = \sigma_1 \sigma_2 \sigma_1$, so that Δ^2 generates the center of \mathbf{B}_3 .

The first step in our algorithm is to turn a word $w \in \mathbf{B}_3$ into its Murasugi normal form. Recall that Murasugi normal forms for conjugacy classes are given by

$$\begin{aligned}
\Omega_0 &= \{\Delta^{2k} \mid k \in \mathbb{Z}\}, \\
\Omega_1 &= \{\Delta^{2k} \sigma_1 \sigma_2 \mid k \in \mathbb{Z}\}, \\
\Omega_2 &= \{\Delta^{2k} (\sigma_1 \sigma_2)^2 \mid k \in \mathbb{Z}\}, \\
\Omega_3 &= \{\Delta^{2k+1} \mid k \in \mathbb{Z}\}, \\
\Omega_4 &= \{\Delta^{2k} \sigma_1^{-p} \mid k, p \in \mathbb{Z}, p > 0\}, \\
\Omega_5 &= \{\Delta^{2k} \sigma_2^q \mid k, q \in \mathbb{Z}, q > 0\}, \\
\Omega_6 &= \{\Delta^{2k} \sigma_1^{-p_1} \sigma_2^{q_1} \cdots \sigma_1^{-p_r} \sigma_2^{q_r} \mid k, r, p_i, q_i \in \mathbb{Z}, r, p_i, q_i > 0\}.
\end{aligned}$$

It is shown in [Mur74, Prop.2.1] that any word $w \in \mathbf{B}_3$ is conjugate to an element of $\Omega_0 \cup \dots \cup \Omega_6$. We first want to ensure that we can obtain this normal form in polynomial time with respect to the length of w . We denote the length of a word w by $|w|$.

Proposition 3.1. *There is an algorithm \mathcal{A}_{Mnf} which turns a word $w \in \mathbf{B}_3$ into its Murasugi normal form $N(w)$, and which runs in linear time with respect to the length of w . Furthermore, the length of $N(w)$ as a word in $\Delta, \sigma_1, \sigma_2$ is linearly bounded in $|w|$.*

Proof. We closely follow the proof of [Mur74, Prop.2.1], turning it into an algorithm with the required properties.

The braid group \mathbf{B}_3 is isomorphic to $G = \langle a, b \mid a^2 = b^3 \rangle$, using the isomorphism $\varphi: G \rightarrow \mathbf{B}_3$ given by $\varphi(a) = \sigma_1 \sigma_2 \sigma_1$ and $\varphi(b) = \sigma_1 \sigma_2$. Then inverse satisfies $\varphi^{-1}(\sigma_1) = b^{-1}a$ and $\varphi^{-1}(\sigma_2) = a^{-1}b^2$. Write $c = a^2 = b^3$ for the element that generates the center of G . Then

$$\varphi^{-1}(\sigma_1) = c^{-1}b^2a, \quad \varphi^{-1}(\sigma_1^{-1}) = c^{-1}ab, \quad \varphi^{-1}(\sigma_2) = c^{-1}ab^2, \quad \varphi^{-1}(\sigma_2^{-1}) = c^{-1}ba,$$

and we get $\varphi^{-1}(w) = c^{-|w|}v_1 \cdots v_{|w|}$ with each $v_i \in \{b^2a, ab, ab^2, ba\}$. The length of $\varphi^{-1}(w)$ in a, b, c is at most $4|w|$. Some of the $v_i v_{i+1}$ can contain a^2 or b^3 in which case we replace them with c and move it to the front. That is, we can write $\varphi^{-1}(w)$ as

$$c^n, c^n a, c^n b^l, \text{ or } c^n a^{\varepsilon_1} b^{l_1} a b^{l_2} \cdots a b^{l_m} a^{\varepsilon_2}, \quad (1)$$

with $n \in \mathbb{Z}$, $l, l_1, \dots, l_m \in \{1, 2\}$, $\varepsilon_1, \varepsilon_2 \in \{0, 1\}$. Furthermore, the length of this word in a, b, c is still at most $4|w|$, and this form can be obtained in linear time depending on $|w|$.

We want to get $\varphi^{-1}(w)$ to be conjugate to one of

$$c^n, c^n a, c^n b^l, \text{ or } c^n a b^{l_1} a b^{l_2} \cdots a b^{l_m}, \quad (2)$$

that is, in the case of the last word in (1) with $\varepsilon_1 = 1$ and $\varepsilon_2 = 0$. Let us check the other cases for this word. If $\varepsilon_1 = 0$ and $\varepsilon_2 = 1$, we have

$$c^n b^{l_1} a b^{l_2} \cdots a b^{l_m} a \sim c^n a b^{l_1} \cdots a b^{l_m},$$

where \sim stands for the conjugation relation. The latter is in the form of (2).

If $\varepsilon_1 = 0 = \varepsilon_2$, we have

$$c^n b^{l_1} a b^{l_2} \cdots a b^{l_m} \sim c^n a b^{l_2} \cdots a b^{l_m + l_1}$$

with $l_m + l_1 \in \{2, 3, 4\}$. If $l_m + l_1 = 2$, this is in the form of (2). If $l_m + l_1 = 4$ we get an extra c , but it is also in the form of (2). If $l_m + l_1 = 3$, we get another c and we have case $\varepsilon_1 = 1 = \varepsilon_2$, but with m replaced by $m - 2$. Note that we can assume $m \geq 2$, since for $m = 1$ the original word is $c^n b^l$. Also, if $m = 2$, we now have the word $c^{n+1} a^2 = c^{n+2}$, which is also in the form of (2).

So, now if $\varepsilon_1 = 1 = \varepsilon_2$, we have

$$c^n a b^{l_1} \cdots a b^{l_m} a \sim c^n a^2 b^{l_1} \cdots a b^{l_m} = c^{n+1} b^{l_1} \cdots a b^{l_m},$$

which brings us back to the case $\varepsilon_1 = 0 = \varepsilon_2$. Nevertheless, after at most m steps, we get $\varphi^{-1}(w)$ to be in the form of (2).

Note that $\varphi(c) = \Delta^2$, $\varphi(ab) = \Delta^2 \sigma_1^{-1}$ and $\varphi(ab^2) = \Delta^2 \sigma_2$, so applying φ to a word in (2) lands in one of $\Omega_0, \dots, \Omega_6$, and this is the required normal form $N(w)$. \square

For words in $\Omega_0, \Omega_1, \Omega_2$ and Ω_3 the integral Khovanov homology of the corresponding braid closure is explicitly given in [CLSS22, Cor.5.7]. For Ω_4 and Ω_5 , the integral Khovanov homology can be obtained using [Kel24, Alg.1]. We will give a slightly different approach using [Sch25b], and express the Khovanov homology of the braid closure of $\Delta^{2k}w \in \Omega_4 \cup \Omega_5 \cup \Omega_6$ in terms of the Khovanov homology of the torus link and the braid closure of w .

Theorem 3.2. *Let L be the braid closure corresponding to the word $\Delta^{2k}\sigma_1^{-l}$ with $k, l \geq 1$. Then*

$$\begin{aligned} \text{Kh}^{4k, 12k-l \pm 1}(L) &= \mathbb{Z} \\ \text{Kh}^{4k-1, 12k-l-1}(L) &= \begin{cases} \mathbb{Z} & l = 1, \\ 0 & l \geq 2, \end{cases} \\ \text{Kh}^{4k-1, 12k-l-3}(L) &= \begin{cases} \mathbb{Z}/2\mathbb{Z} & l = 1, \\ 0 & l \geq 2, \end{cases} \\ \text{Kh}^{4k-2, 12k-l-3}(L) &= \begin{cases} 0 & l = 1, \\ \mathbb{Z}^2 & l = 2, \\ \mathbb{Z} & l \geq 3, \end{cases} \\ \text{Kh}^{4k-2, 12k-l-5}(L) &= \begin{cases} \mathbb{Z} & l = 1, \\ \mathbb{Z}^2 & l = 2, \\ \mathbb{Z}/2\mathbb{Z} & l \geq 3, \end{cases} \\ \text{Kh}^{4k-3, 12k-l-7}(L) &= \begin{cases} 0 & l \leq 2, \\ \mathbb{Z} & l \geq 3, \end{cases} \\ \text{Kh}^{4k-4, 12k-l-7}(L) &= \begin{cases} \mathbb{Z} & l \leq 3, \\ \mathbb{Z}^3 & l = 4, \\ \mathbb{Z}^2 & l \geq 5, \end{cases} \\ \text{Kh}^{4k-4, 12k-l-9}(L) &= \begin{cases} \mathbb{Z} & l \leq 3, \\ \mathbb{Z}^2 & l = 4, \\ \mathbb{Z}/2\mathbb{Z} & l \geq 5. \end{cases} \end{aligned}$$

For any integer j we have

$$\text{Kh}^{4k-5, j}(L) = \text{Kh}^{4k-5, j+l}(T(3, 3k)) \oplus \text{Kh}^{-5, j-12k+1}(T(2, -l)),$$

and for any integers $i \leq 4k-6$ and j we have

$$\text{Kh}^{i, j}(L) = \text{Kh}^{i, j+l}(T(3, 3k)) \oplus \text{Kh}^{i-4k, j-12k+1}(T(2, -l)) \oplus \text{Kh}^{i-4k, j-12k-1}(T(2, -l)).$$

If $k = 1$ we also have

$$\text{Kh}^{0, 12-l-5}(L) = \begin{cases} 0 & l \leq 3, \\ \mathbb{Z} & l \geq 4, \end{cases}$$

and if $k \geq 2$ we also have

$$\begin{aligned} \text{Kh}^{4k-3, 12k-l-3}(L) &= \mathbb{Z} \\ \text{Kh}^{4k-3, 12k-l-5}(L) &= \begin{cases} \mathbb{Z} & l \leq 3, \\ \mathbb{Z}/2\mathbb{Z} & l \geq 4. \end{cases} \end{aligned}$$

In all other bigradings, the homology groups are 0.

We delay the proof of this theorem until the next section.

Theorem 3.3. *Let L be the braid closure corresponding to the word $\Delta^{2k}\sigma_2^l$ with $k, l \geq 1$. For $i < 4k$ we have*

$$\text{Kh}^{i, j}(L) = \text{Kh}^{i, j-l}(T(3, 3k)).$$

For $i > 4k+1$ we have

$$\text{Kh}^{i, j}(L) = \text{Kh}^{i-4k, j-12k+1}(T(2, l)) \oplus \text{Kh}^{i-4k, j-12k-1}(T(2, l)).$$

The other non-zero homology groups are given by

$$\begin{aligned}\mathrm{Kh}^{4k,12k+l-3}(L) &= \mathbb{Z}, \\ \mathrm{Kh}^{4k,12k+l-1}(L) &= \mathbb{Z}^2, \\ \mathrm{Kh}^{4k+1,12k+l+1}(L) &= \mathbb{Z}/2\mathbb{Z}, \\ \mathrm{Kh}^{4k+1,12k+l+3}(L) &= \mathbb{Z}.\end{aligned}$$

This theorem follows directly from [Sch25b, Thm.6.4].

Theorem 3.4. *Let $w = \sigma_1^{-p_1} \sigma_2^{q_1} \cdots \sigma_1^{-p_r} \sigma_2^{q_r}$ with positive integers $r, p_1, q_1, \dots, p_r, q_r$, let L_w be the braid closure corresponding to w and L be the braid closure corresponding to $\Delta^{2k}w$ with $k \geq 1$. Then for integers $i \neq 4k, 4k+1$ and j we have*

$$\mathrm{Kh}^{i,j}(L) = \mathrm{Kh}^{i,j-t}(T(3, 3k)) \oplus \mathrm{Kh}^{i-4k, j-12k}(L_w), \quad (3)$$

with

$$t = q_1 + \cdots + q_r - (p_1 + \cdots + p_r).$$

The remaining non-zero homology groups are given by

$$\begin{aligned}\mathrm{Kh}^{4k,12k+t-1}(L) &= \mathrm{Kh}^{0,t-1}(L_w), \\ \mathrm{Kh}^{4k,12k+t+1}(L) &= \mathrm{Kh}^{0,t+1}(L_w)/\mathbb{Z}, \\ \mathrm{Kh}^{4k+1,12k+t+1}(L) &= \mathrm{Kh}^{1,t+1}(L_w) \oplus \mathbb{Z}/2\mathbb{Z}, \\ \mathrm{Kh}^{4k+1,12k+t+3}(L) &= \mathrm{Kh}^{1,t+3}(L_w) \oplus \mathbb{Z}.\end{aligned}$$

Notice that $\mathrm{Kh}^{0,t+1}(L_w)$ contains at least one direct summand of \mathbb{Z} , and here $\mathrm{Kh}^{0,t+1}(L_w)/\mathbb{Z}$ refers to the homology group with this summand of \mathbb{Z} removed. The proof of this theorem is also delayed until the next section.

In all three cases, the homology behaves as in (3) with a few exceptional bidegrees. Perhaps surprisingly, we have more exceptions in the case corresponding to Ω_4 than to Ω_6 . This is because the closure of the braid word σ_1^{-l} is a split link whose Khovanov homology is not as thin as the Khovanov homology corresponding to the alternating words used in Ω_6 .

Theorems 3.2 - 3.4 assume $k \geq 1$. For negative k , observe that inverting braid words flip the sets Ω_4 and Ω_5 , while the inverse of $\Delta^{2k}\sigma_1^{-p_1}\sigma_2^{q_1} \cdots \sigma_1^{-p_r}\sigma_2^{q_r}$ is conjugate to $\Delta^{-2k}\sigma_1^{-q_r}\sigma_2^{p_r} \cdots \sigma_1^{-q_1}\sigma_2^{p_1}$. We can therefore obtain the Khovanov homology by dualizing the Khovanov homology of the mirror.

The Khovanov homology of $T(3, 3k)$ and $T(2, \pm l)$ is given explicitly in [CLSS22, Cor.5.7] and [Kho00, §6.2], so the only remaining ingredient is the Khovanov homology of the braid closure of $\sigma_1^{-p_1}\sigma_2^{q_1} \cdots \sigma_1^{-p_r}\sigma_2^{q_r}$. There does not seem to be a simple formula for it, but since these links are alternating and non-split, we can derive it from the Jones polynomial and the signature, see [Shu21]. The signature for our braid closure is $t = q_1 + \cdots + q_r - (p_1 + \cdots + p_r)$, compare [Sch25b, Cor.8.3]. For the convenience of the reader, and also to see that deriving the Khovanov homology is done in polynomial time, we now give an explicit formula for the Khovanov homology of a quasi-alternating oriented link L , given the Jones polynomial and the signature $s = s(L)$.

First note that the Jones polynomial is usually given as a Laurent polynomial in a variable $t^{\frac{1}{2}}$. In order to be closer to Khovanov homology we use the substitution $q = t^2$, and we write the Jones polynomial of L as

$$J_L(q) = \sum_{j \in \mathbb{Z}} a_j q^j.$$

It is worth noting that $a_j = 0$ unless $j \in 1 + |L| + 2\mathbb{Z}$, and only finitely many a_j are non-zero. The reduced Khovanov homology groups are easily obtained from the coefficients a_j . Namely, by [MO08] the homology groups are free abelian, with non-zero groups only occurring in bidegrees $(j, s+2j)$, and the Betti numbers given by

$$\tilde{b}^{j,s+2j} = |a_{s+2j}|.$$

Expressing the unreduced Khovanov homology is slightly more difficult. We use the reduced Bar-Natan–Lee–Turner spectral sequence over \mathbb{Z} , which for quasi-alternating links collapses at the E_2 -page, with these groups being free abelian and of total rank $2^{|L|-1}$, compare [Sch25a].

In particular, there exists a homomorphism $d: \tilde{\text{Kh}}^{j,s+2j}(L) \rightarrow \tilde{\text{Kh}}^{j+1,s+2j+2}(L)$ with $d^2 = 0$ and the corresponding homology groups $E_2^{j,s+2j}$ free abelian. Let us denote the Betti numbers for these groups by $\tilde{b}_{\text{BLT}}^{j,s+2j}(L)$. Then $0 \leq \tilde{b}_{\text{BLT}}^{j,s+2j}(L) \leq \tilde{b}^{j,s+2j}(L)$. The exact values can be determined using linking numbers, but for now we simply note that $\tilde{b}_{\text{BLT}}^{j,s+2j}(L) = 0$ for odd j , and $\tilde{b}_{\text{BLT}}^{0,s}(L) \geq 1$. Define

$$\bar{b}^j = \tilde{b}^{j,s+2j}(L) - \tilde{b}_{\text{BLT}}^{j,s+2j}(L).$$

Lemma 3.5. *Let L be a quasi-alternating oriented link with signature s . The Betti numbers of unreduced Khovanov homology are given by*

$$\begin{aligned} b^{j,s+2j-1}(L) &= \sum_{i=j_{\min}}^j (-1)^{j+i} \bar{b}^i + \tilde{b}_{\text{BLT}}^{j,s+2j}(L), \\ b^{j,s+2j+1}(L) &= \sum_{i=j_{\min}}^{j-1} (-1)^{j+i+1} \bar{b}^i + \tilde{b}_{\text{BLT}}^{j,s+2j}(L). \end{aligned}$$

Here $j_{\min} = \min\{j \in \mathbb{Z} \mid \tilde{b}^{j,s+2j} \neq 0\}$. All other Betti numbers are 0. Any torsion elements have order 2, and the only possibly non-zero torsion coefficients are given by

$$t_2^{j,s+2j-1}(L) = \sum_{i=j_{\min}}^{j-1} (-1)^{j+i+1} \bar{b}^i.$$

Proof. We have the long exact sequence

$$0 \longrightarrow \text{Kh}^{j,s+2j+1}(L) \longrightarrow \widetilde{\text{Kh}}^{j,s+2j}(L) \xrightarrow{\delta} \widetilde{\text{Kh}}^{j+1,s+2(j+1)}(L) \longrightarrow \text{Kh}^{j+1,s+2(j+1)-1}(L) \longrightarrow 0,$$

with $\delta = 2d$ by [Sch25a, Lm.5.6]. It follows that

$$b^{j,s+2j-1} = \tilde{b}^{j,s+2j} - \tilde{b}^{j-1,s+2(j-1)} + b^{j-1,s+2(j-1)+1}, \quad (4)$$

and since the reduced Khovanov homology is torsion-free, we have

$$b^{j,s+2j+1} = \text{rank ker}(d: \widetilde{\text{Kh}}^{j,s+2j}(L) \rightarrow \widetilde{\text{Kh}}^{j+1,s+2(j+1)}(L)).$$

Note that we here drop reference to L in the Betti numbers. Let us write d^j to indicate the homological degree of the domain of d . Since $(\widetilde{\text{Kh}}^{j,s+2j}(L), d^j)$ form a cochain complex whose homology is determined by the \tilde{b}_{BLT}^j , we have

$$b^{j,s+2j+1} = \tilde{b}_{\text{BLT}}^j + \text{rank}(d^{j-1}) = \tilde{b}_{\text{BLT}}^j + \tilde{b}^{j-1} - b^{j-1,s+2(j-1)+1}. \quad (5)$$

The statement about the Betti numbers follow via an induction over j , starting with j_{\min} , using (4) and (5). We leave the details to the reader.

For the statement about torsion coefficients, note that $\text{Kh}^{j,s+2j+1}(L)$ is a subgroup of a torsion-free abelian group, hence torsion-free. Since $\delta = 2d$ and the second page of the reduced integral BLT-spectral sequence is torsion-free, the torsion coefficient $t_2^{j,s+2j-1}$ agrees with $\text{rank}(d^{j-1})$. The statement therefore follows from (5). \square

Proof of Theorem 1.3. Let w be a braid word in $\{\sigma_1, \sigma_2, \sigma_1^{-1}, \sigma_2^{-1}\}$. By Proposition 3.1 we get the Murasugi normal form $N(w)$ of w in linear time. If $N(w) \in \Omega_0 \cup \Omega_1 \cup \Omega_2 \cup \Omega_3$, we can read off the Khovanov homology of $L_w = L_{N(w)}$ from [CLSS22, Cor.5.7]. If $N(w) \in \Omega_4 \cup \Omega_5$ with $k \neq 0$ we can read off the Khovanov homology using Theorem 3.2 and 3.3 with [Kho00, §6.2]. If $k = 0$, L_w is a split union of a 2-stranded torus link with an unknot, so we can read off its Khovanov homology using [Kho00, §6.2].

If $N(w) \in \Omega_6$, write $N(w) = \Delta^{2k}v$ with $k \in \mathbb{Z}$ and $v = \sigma_1^{-p_1} \sigma_2^{q_1} \cdots \sigma_1^{-p_r} \sigma_2^{q_r}$. The Jones polynomial of L_v can be obtained in polynomial time, see [MS90], and therefore we can obtain the Khovanov homology of L_v in polynomial time from Lemma 3.5. Note that to compute the $\tilde{b}_{\text{BLT}}^j(L_v)$, we need to work out the number of components of L_v , as well as the differences in linking numbers. Since there can be at most three components, this can be done in linear time. If $k \neq 0$, we use Theorem 3.4 to read off the Khovanov homology of L_w . \square

Algorithm 1 Polynomial time algorithm for Khovanov homology of 3-braids

Input Braid word w on the alphabet $\{\sigma_1, \sigma_2, \sigma_1^{-1}, \sigma_2^{-1}\}$

Output Khovanov homology of the braid closure $\text{Kh}(L_w)$

```

1:  $\Delta^{2k}v \leftarrow$  Murasugi normal form of  $w$ 
2: if  $\Delta^{2k}v \in \Omega_0 \cup \Omega_1 \cup \Omega_2 \cup \Omega_3$  then
3:   Read off  $\text{Kh}(L_w)$  from [CLSS22, Cor.5.7].
4: end if
5: if  $\Delta^{2k}v \in \Omega_4 \cup \Omega_5$  then
6:   Enforce  $k \geq 0$  by mirroring and rotating the braid.
7:   If  $k = 0$ , then  $L_w = T(2, l) \sqcup$  (unknot) and  $\text{Kh}(L_w)$  is obtained with [Kho00, §6.2].
8:   Otherwise, get  $\text{Kh}(L_w)$  with [Kho00, §6.2], [CLSS22, Cor.5.7], Theorem 3.2, and Theorem 3.3.
9: end if
10: if  $\Delta^{2k}v \in \Omega_6$  then
11:   Enforce  $k \geq 0$  by mirroring and rotating the braid.
12:   Compute  $J_L(q)$  where  $L = L_v$  with polynomial time algorithm of [MS90].
13:   Compute the signature of  $L_v$ .
14:   Compute the linking numbers and from them the BLT-generators of  $L_v$ .
15:   Use Lemma 3.5 to combine the above into  $\text{Kh}(L_v)$ 
16:   If  $k = 0$  then  $\text{Kh}(L_w) = \text{Kh}(L_v)$ .
17:   If  $k > 0$  then use  $\text{Kh}(L_v)$ , [CLSS22, Cor.5.7], and Theorem 3.4 to obtain  $\text{Kh}(L_w)$ .
18: end if

```

The algorithm for determining the Jones polynomial of a braid closure given by Morton and Short [MS90] runs in cubic time. Particularly for 3-braids this is very fast. Bar-Natan's algorithm [BN07] on the other hand can be very slow for certain 3-braids. We do not need to look at the details of this algorithm to see why. The main reason is that both Bar-Natan's scanning algorithm and the divide-and-conquer algorithm produce a specific basis for a cochain complex chain homotopy equivalent to the Khovanov complex. Since there are links whose total Betti number grows exponentially, these algorithms spend exponential time on them.

A very simple class of such links is given by 3-stranded *weaving links*. For a positive integer n , the weaving link $W(3, n)$ is the closure of the 3-braid word $(\sigma_1^{-1}\sigma_2)^n$. The determinant of $W(3, n)$ is given by

$$\det(W(3, n)) = \left(\frac{3 + \sqrt{5}}{2}\right)^n + \left(\frac{3 - \sqrt{5}}{2}\right)^n - 2, \quad (6)$$

see [JNP23, Thm.2.2]. Since these weaving links are alternating, the determinant agrees with the total Betti number of reduced Khovanov homology.

For example, the determinant of $W(3, 20)$, a knot with 40 crossings, is 228,826,125. Attempting to run the scanning algorithm on this knot almost certainly leads to a memory overflow, while computing the Jones polynomial is done practically immediately.

Proof of Theorem 1.2. Applying Bar-Natan's scanning algorithm or the divide-and-conquer algorithm to $W(3, n)$ leads to a basis with more than $\det(W(3, n))$ elements, from which the Khovanov homology is determined via standard linear algebra methods. As this basis grows exponentially in n by (6), the number of steps to do this is exponential in the number of crossings of $W(3, n)$.

Now consider the positive word $v = \sigma_1\sigma_2\sigma_1^2\sigma_2^2$. Then

$$v = \sigma_1\sigma_2\sigma_1^2\sigma_2^2 = \sigma_1\sigma_2\sigma_1\sigma_1\sigma_2\sigma_1\sigma_1^{-1}\sigma_2 = \Delta^2\sigma_1^{-1}\sigma_2.$$

In particular, $v^n = \Delta^{2n}(\sigma_1^{-1}\sigma_2)^n$ has $6n$ crossings, and its total Betti number grows exponentially in n by Theorem 3.4. \square

Another family of links with exponentially large Betti number is constructed in Section 7.

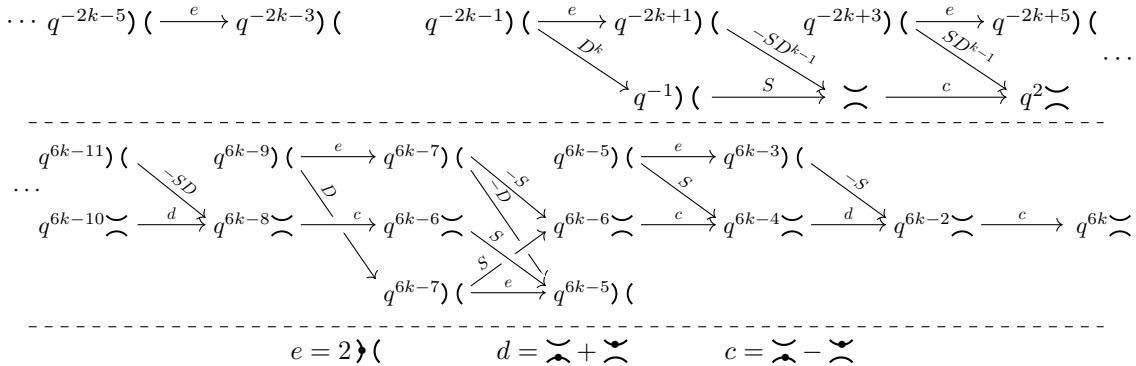
4 Proofs of Theorem 3.2 and 3.4

As with Theorem 3.3, both theorems will follow from results in [Sch25b]. However, we need to fill in more details in these cases. It will be useful to treat links as based links. This way we can view the Khovanov complexes as complexes over $A = \mathbb{Z}[X]/(X^2)$, even though we only view the homology as abelian groups. The base point is placed on the middle strand of the diagram corresponding to the braid word.

4.1 Proof of Theorem 3.2

Let $w = \Delta^{2k} \sigma_1^{-l}$ with $k, l \geq 1$, and let $T_{k,l}$ be the tangle obtained from the braid word w after connecting the two endpoints of the leftmost strand.

Then $q^{l-6k}[[T_{k,l}]]$ is chain homotopy equivalent to a sub-complex of \mathcal{B}_k depicted in Figure 4. More precisely, the sub-complex is obtained by removing all objects from the top row of the form $q^{6k-2r+1}$ (for $r > l$. This follows from [Sch25b, Prop.7.3].



Now consider the functor $G: \text{Cob}_{•/l}^Z(B_2^2) \rightarrow \mathfrak{Mod}_A^q$ obtained by connecting the two endpoints on the right and identifying $\text{Cob}_{•/l}^Z(B_1^1)$ with finitely generated, free q -graded modules over $A = \mathbb{Z}[X]/(X^2)$. Then $G(c) = 0$, so $G(\mathcal{B}_k)$ decomposes into a direct sum of relatively small cochain complexes. These cochain complexes have been analysed in detail in the proof of [Sch25b, Thm.7.4]. In particular, $G(\mathcal{B}_k)$ is chain homotopy equivalent to the complex \mathcal{C}_k given in Figure 5. Moreover, $G(q^{l-6k}[[T_{k,l}]]$) is chain homotopy equivalent to the sub-complex of \mathcal{C}_k obtained by removing all gray objects of homological degree less than $4k - l$.

To prove Theorem 3.2 we only need to analyse the various sub-complexes of \mathcal{C}_k depending on l . First observe that \mathcal{C}_k contains identity morphisms between homological degrees $4k - 2$ and $4k - 1$, between $4k - 3$ and $4k - 2$, and between $4k - 5$ and $4k - 4$ (coming from the $(2X)^{i-1}$ with $i = 1$). To have these identities in the sub-complexes, we additionally need $l \geq 2$ for the first, $l \geq 3$ for the second, and $l \geq 5$ for the third identity. If we can cancel, the morphisms labelled X will disappear, and if we cannot cancel, the morphisms labelled X are not present in the sub-complex.

Furthermore, the diagonal morphisms labelled $(2X)^i$ in the middle part of Figure 5 are only non-zero if $i = 1$, and even if $i = 1$ we can remove them with a change of basis. This implies that up to isomorphism the complex \mathcal{C}_k and its relevant sub-complexes are isomorphic to a direct sum of complexes of the form $u^i q^j A \xrightarrow{2X} u^{i+1} q^{j+2} A$ and $q^j A$. We refer to the former complex as a *knight*, and note its homology is given by $u^i q^{j-1} \mathbb{Z} \oplus u^{i+1} q^{j+1} \mathbb{Z}/2\mathbb{Z} \oplus u^{i+1} q^{j+3} \mathbb{Z}$. The latter complex is called a *pawn*.

The gray knight complexes resemble the Khovanov complexes of the split union of $T(2, -l)$ with an unknot, while the black knight and pawn complexes give the Khovanov homology of $T(3, 3k)$, except

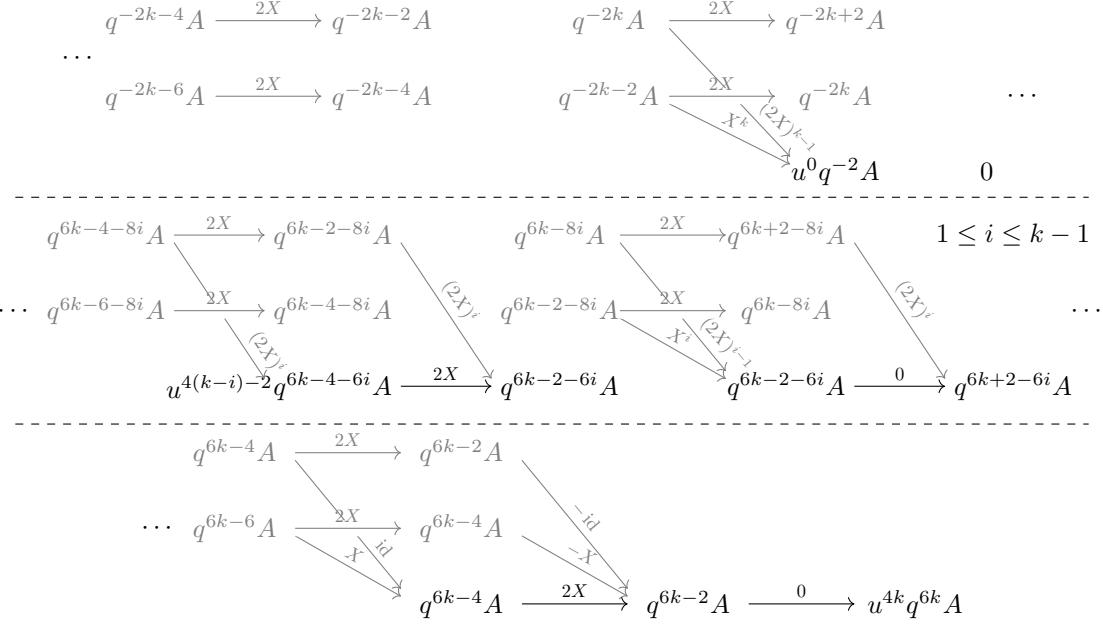


Figure 5: The cochain complex \mathcal{C}_k . Notice that $(2X)^i$ is 0 for $i \geq 2$, and 1 for $i = 0$. Homological degrees are indicated by u^i for some modules.

in homological degrees near $4k$. Therefore we get in homological degrees $i \leq 4k - 6$ the direct sums of Khovanov homologies for $T(3, 3k)$ and the split union of $T(2, l)$ with an unknot up to grading shifts.

For the homological degrees between $4k - 5$ and $4k$ we need to check various cases, since there is some interaction between the gray and black parts. In particular, the value of l determines which cancellations we can do. In addition to distinguishing between $k = 1$ and $k \geq 2$, we need to consider $l \in \{1, 2, 3, 4\}$ and $l \geq 5$. While there are several cases to check, they are straightforward and left to the reader.

4.2 Proof of Theorem 3.4

Let $w = \sigma_1^{-p_1} \sigma_2^{q_1} \cdots \sigma_1^{-p_r} \sigma_2^{q_r}$ with positive integers $r, p_1, q_1, \dots, p_r, q_r$. By [Sch25b, Thm.9.2] we have that

$$\text{CKh}(L_{\Delta^{2k}w}) \simeq B_k \oplus D_w,$$

where B_k is a direct sum of knight and pawn complexes concentrated in homological degrees less than $4k - 4$, whose homology satisfies

$$H^{i,j-t}(B_k) \cong \text{Kh}^{i,j}(T(3, 3k))$$

for $i \leq 4k - 5$ and $t = q_1 + \cdots + q_r - (p_1 + \cdots + p_r)$, and D_w fits into a short exact sequence of A -cochain complexes

$$0 \longrightarrow T \longrightarrow D_w \longrightarrow u^{4k}q^{12k}C_w \longrightarrow 0,$$

where $C_w \oplus q^t A \simeq \text{CKh}(L_w)$ as A -cochain complexes, and

$$T \cong \begin{cases} u^{4k-4}q^{12k+t-8}A \oplus u^{4k-3}q^{12k+t-4}A \oplus u^{4k-2}q^{12k+t-4}K \oplus u^{4k}q^{12k+t}K & k \geq 2 \\ q^{t+4}A \oplus u^2q^{t+8}K \oplus u^4q^{t+12}K & k = 1 \end{cases}$$

Here $K = (A \xrightarrow{2X} uq^2A)$ is a knight complex. From this it follows directly that

$$\text{Kh}^{i,j}(L_{\Delta^{2k}w}) \cong \text{Kh}^{i,j-t}(T(3, 3k)) \oplus \text{Kh}^{i-4k, j-12k}(L_w) \quad (7)$$

for $i \leq 4k - 5$ and $i \geq 4k + 2$. It remains to check that this also holds for $i \leq 4k - 1$, and for $i \in \{4k, 4k + 1\}$ we get the formulas as in the statement of Theorem 3.4.

Let us write $s = 12k + t$, and consider the long exact sequence

$$\cdots \longrightarrow \tilde{H}^{4k+i-1, s+2i-2}(D_w) \xrightarrow{\delta} \tilde{H}^{4k+i, s+2i}(D_w) \longrightarrow H^{4k+i, s+2i-1}(D_w) \longrightarrow \tilde{H}^{4k+i, s+2i-2}(D_w) \quad (8)$$

Here, \tilde{H} refers to the homology of the complex $D_w \otimes_A \mathbb{Z}$, where X acts as 0 on \mathbb{Z} . For $i \geq -5$ we have

$$\begin{aligned} \tilde{H}^{4k+i, s+2i}(D_w) &\cong \widetilde{\text{Kh}}^{4k+i, s+2i}(L_{\Delta^{2k} w}) \cong \tilde{H}^{i, t+2i}(C_w) \oplus \tilde{H}^{4k+i, s+2i}(T) \\ \tilde{H}^{4k+i, s+2i-2}(D_w) &\cong \widetilde{\text{Kh}}^{4k+i, s+2i-2}(L_{\Delta^{2k} w}) \cong 0, \end{aligned}$$

with $\tilde{H}^{i,j}(C_w) = \widetilde{\text{Kh}}^{i,j}(L_w)$ for $(i, j) \neq (0, t)$, and $H^{0,t}(C_w) \oplus \mathbb{Z} \cong \widetilde{\text{Kh}}^{0,t}(L_w)$ by [Sch25b, Cor.9.3]. Furthermore, δ is the connecting homomorphism of the reduced/unreduced long exact sequence in Khovanov homology of $L_{\Delta^{2k} w}$, which by [Sch25a, Lm.5.6] is twice the differential of the first page of the reduced integral BLT spectral sequence.

Notice that $\tilde{H}^{4k+i, s+2i}(T) \cong \mathbb{Z}$ for $i = -4, -2, -1, 0, 1$. We also have $\tilde{H}^{4k-3, s-6}(T) = 0$, and if $k \geq 2$ we have

$$\tilde{H}^{4k-3, s-4}(T) \cong \mathbb{Z} \cong \widetilde{\text{Kh}}^{4k-3, s-4}(L_{\Delta^{2k} w}).$$

In the reduced integral BLT spectral sequence the differential on the first page induces isomorphisms $\tilde{H}^{4k+i, s+2i}(T) \cong \tilde{H}^{4k+i+1, s+2i+2}(T)$ for $i = -2, 0$, while $\mathbb{Z} \cong \tilde{H}^{4k-4, s-8}(T)$ survives to the E_2 -page. The long exact sequence (8) with $i = -4$ then turns into

$$\widetilde{\text{Kh}}^{-5, t-10}(L_w) \xrightarrow{\delta} \widetilde{\text{Kh}}^{-4, t-8}(L_w) \oplus \mathbb{Z} \longrightarrow \text{Kh}^{4k-4, s-9}(L_{\Delta^{2k} w}) \longrightarrow 0$$

and the summand of \mathbb{Z} injects into $\text{Kh}^{4k-4, s-9}(L_{\Delta^{2k} w})$. In particular,

$$\text{Kh}^{4k-4, s-9}(L_{\Delta^{2k} w}) \cong \text{Kh}^{-4, t-9}(L_w) \oplus \mathbb{Z} \cong \text{Kh}^{-4, t-9}(L_w) \oplus \text{Kh}^{4k-4, 12k-9}(T(3, 3k)).$$

The exact same argument works for $i = -2$, thus establishing (7) for $(4k-2, s-5)$, and for $i = 0$ we get

$$\text{Kh}^{4k, s-1}(L_{\Delta^{2k} w}) \cong H^{0, t-1}(C_w) \oplus \mathbb{Z},$$

but as was noted above, $H^{0, t-1}(C_w) \oplus \mathbb{Z} \cong \text{Kh}^{0, t-1}(L_w)$.

For $i = -3$ the long exact sequence (8) is

$$\widetilde{\text{Kh}}^{-4, t-8}(L_w) \oplus \mathbb{Z} \xrightarrow{\delta} \widetilde{\text{Kh}}^{-3, t-6}(L_w) \longrightarrow \text{Kh}^{4k-3, s-7}(L_{\Delta^{2k} w}) \longrightarrow 0,$$

with the \mathbb{Z} -summand in the kernel of δ . In particular, (7) holds for $(i, j) = (4k-3, s-7)$ as the torus link does not contribute to this bigrading.

For $i = -1$ the sequence (8) is

$$\begin{array}{ccccccc} \widetilde{\text{Kh}}^{4k-2, s-4}(L_{\Delta^{2k} w}) & \xrightarrow{\delta} & \widetilde{\text{Kh}}^{4k-1, s-2}(L_{\Delta^{2k} w}) & \longrightarrow & \text{Kh}^{4k-1, s-3}(L_{\Delta^{2k} w}) & \longrightarrow & 0 \\ \downarrow \cong & & \downarrow \cong & & \downarrow \cong & & \\ \widetilde{\text{Kh}}^{-2, t-4}(L_w) \oplus \mathbb{Z} & \xrightarrow{\delta} & \widetilde{\text{Kh}}^{-1, t-2}(L_w) \oplus \mathbb{Z} & \longrightarrow & \text{Kh}^{4k-1, s-3}(L_{\Delta^{2k} w}) & \longrightarrow & 0 \end{array}$$

Recall that δ is twice the differential of the first page of the reduced integral BLT-spectral sequence, which cancels the two \mathbb{Z} -summands. It follows that

$$\text{Kh}^{4k-1, s-3}(L_{\Delta^{2k} w}) \cong \text{Kh}^{-1, t-3}(L_w) \oplus \mathbb{Z}/2\mathbb{Z},$$

with $\mathbb{Z}/2\mathbb{Z} \cong \text{Kh}^{4k-1, 12k-3}(T(3, 3k))$.

The case $i = 1$ is very similar. The main difference is that $\widetilde{\text{Kh}}^{4k, s}(L_{\Delta^{2k} w}) \cong \tilde{H}^{0, t}(C_w) \oplus \mathbb{Z}$. Recall that $\tilde{H}^{0, t}(C_w) \oplus \mathbb{Z} \cong \widetilde{\text{Kh}}^{0, t}(L_w)$, with this summand of \mathbb{Z} in the kernel of δ . We thus also get

$$\text{Kh}^{4k+1, s+1}(L_{\Delta^{2k} w}) \cong \text{Kh}^{1, t+1}(L_w) \oplus \mathbb{Z}/2\mathbb{Z},$$

which is one of the exceptional cases in Theorem 3.4. It remains to prove the theorem in bigradings $(4k+i, s+2i+1)$ for $i = -4, \dots, 1$ and $(4k-3, s-3)$.

In bidegree $(4k-3, s-3)$ only the torus link contributes, and it is \mathbb{Z} if $k \geq 2$ and 0 if $k=1$. But this is exactly the contribution of T in this bidegree.

Now consider the long exact sequence

$$\cdots \longrightarrow \tilde{H}^{4k+i, s+2i+2}(D_w) \longrightarrow H^{4k+i, s+2i+1}(D_w) \longrightarrow \tilde{H}^{4k+i, s+2i}(D_w) \xrightarrow{\delta} \tilde{H}^{4k+i+1, s+2i+2}(D_w) \quad (9)$$

The group $\tilde{H}^{4k+i, s+2i+2}(D_w) = 0$ unless $i = -3$ and $k \geq 2$, in which case it is \mathbb{Z} that injects into $H^{4k-3, s-5}(D_w)$. Then

$$\tilde{H}^{4k-3, s-6}(D_w) \cong \widetilde{\text{Kh}}^{4k-3, s-6}(L_{\Delta^{2k} w}) \cong \widetilde{\text{Kh}}^{-3, t-6}(L_w).$$

The kernel of δ is the kernel of the first differential in the reduced integral BLT-spectral sequence. Hence

$$\text{Kh}^{4k-3, s-5}(L_{\Delta^{2k} w}) \cong \mathbb{Z} \oplus \text{Kh}^{-3, t-5}(L_w),$$

establishing (7) for $(4k-3, s-5)$. Notice that we always have the long exact sequence

$$0 \longrightarrow \text{Kh}^{i, t+2i+1}(L_w) \longrightarrow \widetilde{\text{Kh}}^{i, t+2i}(L_w) \xrightarrow{\delta} \widetilde{\text{Kh}}^{i+1, t+2i+2}(L_w) \longrightarrow \cdots$$

For the other relevant values of i we have

$$\text{Kh}^{4k+i, s+2i+1}(L_{\Delta^{2k} w}) \cong \ker(\delta: \tilde{H}^{4k+i, s+2i}(D_w) \rightarrow \tilde{H}^{4k+i+1, s+2i+2}(D_w)).$$

For $i = -4$ we have $\tilde{H}^{4k-4, s-8}(D_w) \cong \widetilde{\text{Kh}}^{-4, t-8}(L_w) \oplus \mathbb{Z}$, with $\mathbb{Z} \cong \widetilde{\text{Kh}}^{4k-4, 12k-8}(T(3, 3k))$ in $\ker \delta$. Hence

$$\text{Kh}^{4k-4, s-7}(L_{\Delta^{2k} w}) \cong \text{Kh}^{-4, t-7}(L_w) \oplus \mathbb{Z} \cong \text{Kh}^{-4, t-7}(L_w) \oplus \text{Kh}^{4k-4, 12k-8}(T(3, 3k)).$$

The same argument applies to $i = -1$, establishing (7) for $(4k-1, s-1)$.

For $i = -2$ we have

$$\tilde{H}^{4k-2, s-4}(D_w) \cong \widetilde{\text{Kh}}^{-2, t-4}(L_w) \oplus \mathbb{Z}.$$

This time $\mathbb{Z} \cong \widetilde{\text{Kh}}^{4k-2, 12k-4}(T(3, 3k))$ is not in $\ker \delta$. But since $\text{Kh}^{4k-2, 12k-3}(T(3, 3k)) = 0$ we still get (7) in bidegree $(4k-2, s-3)$.

For $i = 0$ we have

$$\tilde{H}^{4k, s}(D_w) \cong \tilde{H}^{0, t}(C_w) \oplus \mathbb{Z} \cong \widetilde{\text{Kh}}^{0, t}(L_w),$$

but δ restricted to the summand \mathbb{Z} is not 0 (recall that $\widetilde{\text{Kh}}^{0, t}(L_w)$ contains a summand \mathbb{Z} which survives the reduced BLT-spectral sequence, but we have a different behaviour here). It follows that

$$\text{Kh}^{4k, s+1}(L_{\Delta^{2k} w}) \cong \text{Kh}^{0, t+1}(L_w)/\mathbb{Z},$$

as stated in Theorem 3.4.

For the remaining case $i = 1$ observe that $\tilde{H}^{4k+1, s+2}(D_w)$ has an extra summand \mathbb{Z} coming from the homology of T , which produces an extra generator in $\ker \delta$. Therefore

$$\text{Kh}^{4k+1, s+3}(L_{\Delta^{2k} w}) \cong \text{Kh}^{1, t+3}(L_w) \oplus \mathbb{Z},$$

as claimed.

5 Extremal Khovanov homology of braids in polynomial time

In this section we prove Theorem 1.4 by constructing algorithms \mathcal{A}_k which work on link diagrams and can be specialized into braid algorithms $\mathcal{B}_{k,t}$. We start in (5.1) by giving an overview of algorithms \mathcal{A}_k constructing $\mathcal{B}_{k,t}$ using them. In (5.2) the algorithms \mathcal{A}_k are described in more detail and in (5.3) and (5.4) their validity and polynomial runtime are proven.

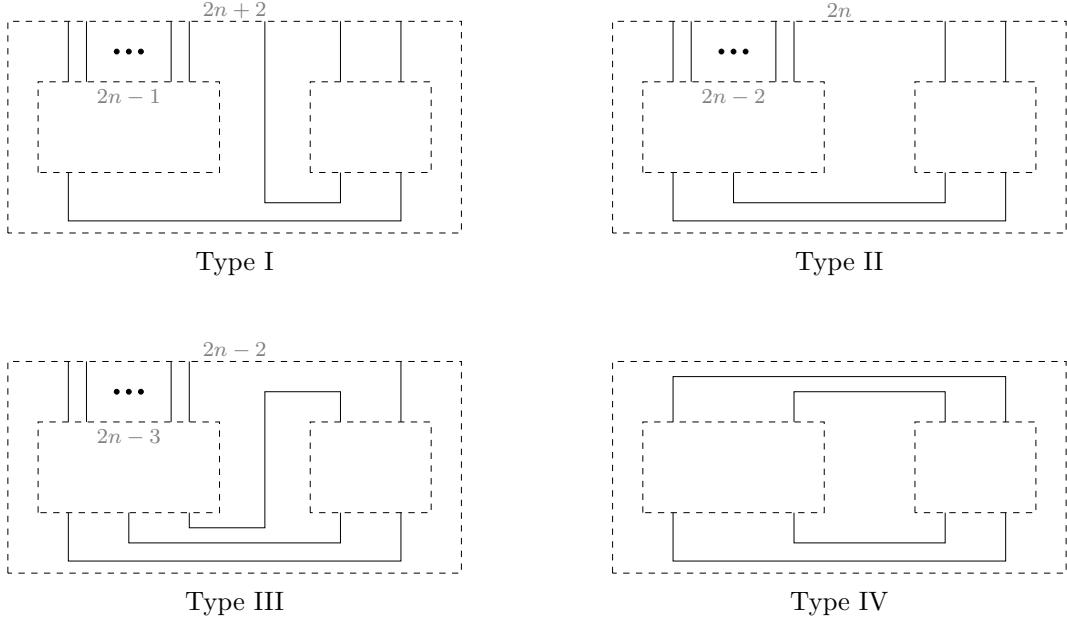


Figure 6: The four different families of simple planar arc diagrams which we will use. Type I, II and III simple planar arc diagrams are determined by $n \geq 1$ and Type IV is unique. The planar arc diagrams are considered up to boundary preserving planar isotopy.

5.1 Overview of the algorithm \mathcal{A}_k

Planar arc diagrams glue “smaller” tangles together to generate “larger” tangles and eventually links. Bar-Natan’s theory allows one to formulate local Khovanov complexes of “smaller” tangles, which themselves can be composed as well with planar arc diagrams and the two ways of making a “larger” composite complex commute with each other [BN05]. Instead of using the general theory, it suffices for our purposes to introduce unoriented *simple planar arc diagrams* of Type I, II, III, and IV, see Figure 6.

A *nice scanning sequence* \mathcal{S} of a link diagram L with n crossings is a triple of sequences

$$\mathcal{S} = ((T_1, \dots, T_n), (c_1, \dots, c_n), (P_1, \dots, P_{n-1}))$$

where T_i are unoriented subtangle diagrams of L , c_i are unoriented crossings of L and P_i are simple planar arc diagrams of Type I, II or III for $i \leq n-2$ and P_{n-1} is of type IV. Additionally, we require that $T_1 = c_1$, $T_n = L$ and $T_i \otimes_{P_i} c_{i+1} = T_{i+1}$ as tangle/link diagrams. A scanning sequence is also implicitly present in the scanning algorithm of [BN07], but our technical condition of allowing only one Type IV diagram is the extra ingredient which makes our sequences *nice*. In Lemma 6.2 we prove that for reasonable link diagrams nice scanning sequences always exist. The girth of a nice scanning sequence, denoted by $\text{girth}(\mathcal{S})$, is the maximum number of boundary points on all subtangles T_i .

Proposition 5.1. *For every $k \geq 0$ there is an algorithm \mathcal{A}_k which takes in a link diagram L and a nice scanning sequence \mathcal{S} for L and outputs $\text{Kh}^{i,*}(L)$ for $i < -n_-(L) + k$ and $i > n_+(L) - k$. The algorithm \mathcal{A}_k has running time $\mathcal{O}(f(\text{girth}(\mathcal{S})) \cdot p_k(n(L)))$ for some function f and polynomial p_k .*

Theorem 1.4 is obtained as a corollary Proposition 5.1.

Proof of Theorem 1.4. The algorithm $\mathcal{B}_{k,t}$ takes in a word w applies the algorithm \mathcal{A}_{k+t-1} to the pair (L_b, \mathcal{S}) . The link diagram L_b is the braid closure of $b = \sigma_1 \dots \sigma_{t-1} w \sigma_{t-1}^{-1} \dots \sigma_1^{-1}$ and \mathcal{S} scans the crossings of L_b in the order that they occur in b . Conjugating w with $\sigma_1 \dots \sigma_{t-1}$ before taking the braid closure does not alter the isotopy type of the link but it will ensure that \mathcal{S} is a nice scanning sequence with $\text{girth}(\mathcal{S}) = 2t$. The conjugation also adds $t-1$ positive and negative crossings, which is why we need to apply \mathcal{A}_{k+t-1} to L_b in order to obtain $\text{Kh}^{i,*}(L_w) = \text{Kh}^{i,*}(L_b)$ in homological degrees $i \leq -n_-(L_w) + k = -n_-(L_b) + k + t - 1$ and $i \geq n_+(L_w) - k = -n_+(L_b) - k - t + 1$. \square

To prove Proposition 5.1 it suffices to describe an algorithm which obtains $\text{Kh}^{i,*}(L)$ for $i < -n_-(L) + k$. The degrees $i > n_+(L) - k$ can be obtained by applying the same algorithm to the mirrored diagram $D^!$ and using the fact that Khovanov homology of the mirror link is obtained from the dual complex. In order to minimize the trouble of grading shifts, we will make use of unshifted complexes and delay the shifts to the very end. Thus, for a crossing $c \in \{\times, \times\}$ we denote $\llbracket c \rrbracket_{\text{US}}$ as the *unshifted Bar-Natan complex* of c , so that

$$\begin{aligned}\llbracket \times \rrbracket_{\text{US}} &= \cdots \rightarrow 0 \rightarrow u^0 q^0) \xrightarrow{\text{saddle}} u^1 q^1 \asymp \rightarrow 0 \rightarrow \cdots \\ \llbracket \times \rrbracket_{\text{US}} &= \cdots \rightarrow 0 \rightarrow u^0 q^0 \asymp \xrightarrow{\text{saddle}} u^1 q^1) \rightarrow 0 \rightarrow \cdots.\end{aligned}$$

and more generally for a tangle diagram T we write $\llbracket T \rrbracket_{\text{US}} = u^{n_-(L)} q^{-n_+(L) + 2n_-(L)} \llbracket T \rrbracket$.

The algorithms \mathcal{A}_k are roughly described by pseudocode in Algorithm 2 and in Subsection 5.2, we will express the precise meaning of Lines 4 and 5 whose point that we do not eliminate “too many” isomorphisms. Limiting the number of cancellations allows us to keep the lengths of integers in check, which ensures that we do not spend exponential time multiplying integers of superpolynomial bit-length.

Algorithm 2 For any k , the following describes \mathcal{A}_k for the lowest degrees and with \mathbb{Z} -coefficients.

Input Link diagram L with a nice scanning sequence $((T_1, \dots, T_{n(L)}), (c_1, \dots, c_{n(L)}), (P_1, \dots, P_{n(L)-1}))$.

Output $\text{Kh}^{i,*}(L)$ for $i \leq k - n_-(L)$

- 1: Assign $C \leftarrow \llbracket c_1 \rrbracket_{\text{US}}$
 - 2: **for** $t = 1, \dots, n(L) - 1$ **do**
 - 3: Assign $C' \leftarrow C \otimes_{P_t} \llbracket c_{t+1} \rrbracket_{\text{US}}$
 - 4: Deloop the newly formed circles in C' and fix a block decomposition. Call the result C'' .
 - 5: Gaussian eliminate certain blocks of C'' by using the saddle of $\llbracket c_{j+1} \rrbracket_{\text{US}}$. Call the result C''' .
 - 6: Truncate C''' at $k + 1$, that is, set $(C''')^{k+2} = 0$. Call the result C .
 - 7: **end for**
 - 8: On C replace every \emptyset with \mathbb{Z} . (Or to put it fancily, apply the Khovanov TQFT functor F .)
 - 9: Put the differentials of $F(C)$ into Smith normal form.
 - 10: Read of the homology from the Smith normal form and shift it with $u^{-n_-(L)} q^{n_+(L) - 2n_-(L)}$.
-

Over finite fields, one does not need to restrain the number of cancellations and pseudocode in Algorithm 3 presents simpler algorithms $\mathbb{F}_p \mathcal{A}_k$ which compute the homology with \mathbb{F}_p coefficients. Removing Line 6 from the Algorithms $\mathbb{F}_p \mathcal{A}_k$ gives precisely Bar-Natan’s scanning algorithm [BN07] which computes the whole Khovanov homology table $\text{Kh}^{*,*}(L; \mathbb{F}_p)$.

Algorithm 3 For any k the following describes $\mathbb{F}_p \mathcal{A}_k$ for the lowest degrees and with \mathbb{F}_p coefficients.

Input Link diagram L with a nice scanning sequence $((T_1, \dots, T_{n(L)}), (c_1, \dots, c_{n(L)}), (P_1, \dots, P_{n(L)-1}))$.

Output $\text{Kh}^{i,*}(L; \mathbb{F}_p)$ for $i \leq k - n_-(L)$

- 1: Assign $C \leftarrow \llbracket c_1 \rrbracket_{\text{US}}$
 - 2: **for** $t = 1, \dots, n(L) - 1$ **do**
 - 3: Assign $C' \leftarrow C \otimes_{P_t} \llbracket c_{t+1} \rrbracket_{\text{US}}$
 - 4: Deloop the newly formed circles in C' . Call the result C''
 - 5: Iteratively Gaussian eliminate isomorphisms of C'' until none are left. Call the result C''' .
 - 6: Truncate C''' at $k + 1$, that is, set $(C''')^{k+2} = 0$. Call the result C .
 - 7: **end for**
 - 8: Replace every \emptyset with \mathbb{Z} . (Or to put it fancily, apply the Khovanov TQFT functor F .)
 - 9: Read of the homology from C and shift it with $u^{-n_-(L)} q^{n_+(L) - 2n_-(L)}$.
-

5.2 Expanding and eliminating blocks in \mathcal{A}_k

Every object B in $\text{Cob}_{\bullet/l}^{\mathbb{Z}}(2n)$ is isomorphic to a unique direct sum (up to a permutation)

$$B \cong \bigoplus_{i=1}^n q^{c_i} D_i \tag{10}$$

where D_i are loopless diagrams, which are planar matchings on $2n$ points, and $c_i \in \mathbb{Z}$. Using (10) any morphism $f: A \rightarrow B$ of $\mathcal{Cob}_{\bullet/l}^{\mathbb{Z}}(2n)$ can be described with a matrix, where each matrix element is a morphism $f_{i,j} \in \text{Hom}(A_j, B_i)$ between some loopless diagrams A_j and B_i .

A morphism $g: D \rightarrow E$ between loopless diagrams is a formal sum of oriented surfaces embedded in $[0, 1]^3$ with boundary $X = (D \times \{1\}) \cup (E \times \{0\}) \cup (\partial D \times [0, 1])$. Topologically, X is a disjoint union of \mathbb{S}^1 :s and we can glue disks to each loop and embed them inside the cube $[0, 1]^3$. Up to a boundary preserving diffeomorphism, this generates a unique cobordism $\Sigma_{D,E} \hookrightarrow [0, 1]^3$. We define $\text{dec}(D, E)$ to be the set of all possible ways of decorating each component of $\Sigma_{D,E}$ with 0 or 1 dots. In other words, $\text{dec}(D, E)$ consists of 2^c copies of $\Sigma_{D,E}$ with varying decorations where c is the number of \mathbb{S}^1 :s in X .

Lemma 5.2. *For loopless planar matchings $D, E \in \mathcal{Cob}_{\bullet/l}^{\mathbb{Z}}(2n)$, the set $\text{dec}(D, E)$ freely generates $\text{Hom}(D, E)$.*

Proof. To see that $\text{dec}(D, E)$ generates $\text{Hom}(D, E)$, take any dotted surface S embedded $[0, 1]^3$ with the desired boundary. By using the neck-cutting relation, one can get rid of all of the genus and disconnect every boundary component of S . Then S is expressed as a sum of unions of dotted spheres and disks. Using the double dot and sphere relations to this sum, we get a description of S as a sum over $\text{dec}(D, E)$.

Next, let $\sum_{S \in \text{dec}(D, E)} c_S S = \sum_{S \in \text{dec}(D, E)} d_S S$ for some integers c_S and d_S . To see that $c_S = d_S$ for a fixed S , we describe a map $f_S: \text{Hom}(D, E) \rightarrow \text{Hom}(\emptyset, \emptyset)$. The map f_S takes a dotted surface Σ in $[0, 1]^3$ whose boundary is a disjoint union of \mathbb{S}^1 :s in $\partial([0, 1]^3)$ and glues disks to each of those boundary loops. The glued discs are embedded “outside of the box” and f_S places a dot on each “outside disk” if the corresponding “inside disk” has no dot in the surface S . Rescaling the generated surface appropriately, we get a surface $f_S(\Sigma) \in \text{Hom}(\emptyset, \emptyset)$. Moreover, this procedure describes a well-defined linear map $f_S: \text{Hom}(D, E) \rightarrow \text{Hom}(\emptyset, \emptyset)$. By plugging in the original sums and using the sphere and double-dot relations, we obtain

$$c_S E = f_S \left(\sum_{S' \in \text{dec}(D, E)} c_{S'} S' \right) = f_S \left(\sum_{S' \in \text{dec}(D, E)} d_{S'} S' \right) = d_S E$$

where E is the empty cobordism $E: \emptyset \rightarrow \emptyset$. Since $\text{Hom}(\emptyset, \emptyset)$ is freely generated by E , this concludes the proof. \square

In complexity theory, finite fields are often simple to work with since addition and multiplication of two elements can be done in constant time in terms of bit-operations on a computer. Storing an integer c takes $\log_2 |c|$ amount of memory and time which can be too much if c is humongous. Asymptotically optimal algorithms for multiplying integers have been extensively studied in computer science, but for us it will suffice that both addition $c + d$ and multiplication cd can be computed in polynomial time with respect to $\log_2(|cd|)$. Nevertheless, in order to prove Proposition 5.1 we will need to make sure that the $\log_2(|c|)$ is bounded by a polynomial for all of the integers involved. To this end, we define $\|\cdot\|$ to be the l_1 norm on $\text{Hom}(D, E)$ with respect to the basis $\text{dec}(D, E)$, that is,

$$\left\| \sum_{S \in \text{dec}(D, E)} c_S S \right\| = \sum_S |c_S|.$$

Lemma 5.3. *There is an algorithm which takes in loopless planar matchings $A, B, C \in \mathcal{Cob}_{\bullet/l}^{\mathbb{Z}}(2n)$ and morphisms $f_1: A \rightarrow B$, $f_2: B \rightarrow C$ expressed in the basis as*

$$f_1 = \sum_{S \in \text{dec}(A, B)} c_S S, \quad f_2 = \sum_{S' \in \text{dec}(B, C)} d_{S'} S'$$

and outputs

$$f_2 f_1 = \sum_{S'' \in \text{dec}(A, C)} e_{S''} S''.$$

For a fixed n , the algorithm runs in polynomial time with respect to $\log(\|f_2\| \cdot \|f_1\|)$ and the composition satisfies $\|f_2 f_1\| \leq 2^{n-1} \|f_2\| \cdot \|f_1\|$.

Proof. Our first goal is to isolate the decorations and coefficients of f_1 and write

$$f_1 = \sum_{S \in \text{dec}(A, B)} c_S S = \Sigma_{A, B} \circ \sum_{P \in \text{dec}(A, A)} t_P P \quad (11)$$

where the underlying surface $\Sigma_{A, A}$ of $\text{dec}(A, A)$ is the product cobordism $A \times [0, 1]$. On each dot decoration $S \in \text{dec}(A, B)$ one has to choose to which components of $A \times [0, 1]$ the dots are moved, that is, one has to pick an appropriate decoration $P \in \text{dec}(A, A)$. For these choices $t_P = c_S$ and the rest of the coefficients t_P vanish. This quick procedure (non-uniquely) finds (11) and similarly expanding f_2 we can obtain

$$f_2 f_1 = \sum_{Q \in \text{dec}(C, C)} r_Q Q \circ \Sigma_{B, C} \circ \Sigma_{A, B} \circ \sum_{P \in \text{dec}(A, A)} t_P P. \quad (12)$$

Next, we will analyze the non-decorated surface $\Sigma_{B, C} \Sigma_{A, B}$. The connected components $\{\Sigma_i\}_i$ of $\Sigma_{B, C} \Sigma_{A, B}$ can be obtained directly from the planar matchings A, B , and C . Each connected Σ_i is obtained by gluing k discs of $\Sigma_{A, B}$ and l discs of $\Sigma_{B, C}$ along m lines yielding $\chi(\Sigma_i) = k + l - m$. By the classification of oriented surfaces with boundary we also get $\chi(\Sigma_i) = 2 - r - 2g$ where r is the number of boundary loops of Σ_i and g is its genus. Since k, l, m and r are obtainable from A, B and C we can compute g . By using the neck cutting relation, we can get rid of the genus on Σ_i at the cost of adding g dots and a coefficient 2^g . Further using the neck-cutting relation $r - 1$ times we can replace all connections of Σ_i with a sum of decorated discs which are glued on the r loops. Applying this simplification to every connected component Σ_i gives the expansion

$$\Sigma_{B, C} \Sigma_{A, B} = \sum_{S'' \in \text{dec}(A, C)} b_{S''} S''. \quad (13)$$

Finally, the expansion of $f_2 f_1$ can be computed from (12) by adding and multiplying the integers t_P , r_Q , $b_{S''}$ and using the double dot relation.

Since $\Sigma_{A, B}$ and $\Sigma_{B, C}$ contain at least 1 disk each and $\Sigma_{B, C} \Sigma_{A, B}$ is obtained by gluing them along n lines, we get $\chi(\Sigma_{B, C} \Sigma_{A, B}) \geq 2 - n$. On the other hand $\Sigma_{A, C}$ contains at most n disks, so $\chi(\Sigma_{A, C}) \leq n$. Thus $\chi(\Sigma_{A, C}) - \chi(\Sigma_{B, C} \Sigma_{A, B}) \leq 2n - 2$ which means that the neck cutting relation can be used at most $n - 1$ times when arriving at (13). Hence $\|\Sigma_{B, C} \Sigma_{A, B}\| \leq 2^{n-1}$ and we can compute

$$\begin{aligned} \|f_2 f_1\| &= \left\| \sum_{Q \in \text{dec}(C, C)} r_Q Q \circ \Sigma_{B, C} \circ \Sigma_{A, B} \circ \sum_{P \in \text{dec}(A, A)} t_P P \right\| \\ &\leq \left\| \sum_{Q \in \text{dec}(C, C)} r_Q Q \right\| \cdot \|\Sigma_{B, C} \Sigma_{A, B}\| \cdot \left\| \sum_{P \in \text{dec}(A, A)} t_P P \right\| \\ &\leq 2^{n-1} \|f_2\| \cdot \|f_1\| \end{aligned}$$

concluding the proof. \square

Now we can elaborate on what we mean by delooping, block decompositions and cancellations at Lines 4 and 5 in Algorithm 2. At Line 4 the complex C'' with a block decomposition and a Morse matching M are formed. Then at Line 5 the Morse complex $M(C'')$ is computed and renamed to be C''' .¹ The concrete block decomposition of C'' and the matching M depend on the type of the simple planar arc diagram P_j , denoted by P for simplicity. **Type I** is easy: no new loops are formed, we take C'' to have the trivial block decomposition and set $M = \emptyset$. Simply put, if P is a Type I planar arc diagram Lines 4 and 5 do nothing.

Type II: By definition, the complex $C' = C \otimes_P \llbracket c \rrbracket_{\text{US}}$ can be written as

$$\begin{array}{ccccccc} C^i \otimes_P & \xrightarrow{\quad} & C^{i+1} \otimes_P & \xrightarrow{\quad} & C^{i+2} \otimes_P & \xrightarrow{\quad} & \dots \\ \dots & & \searrow (-1)^i S^i & & \searrow (-1)^{i+1} S^{i+1} & & \dots \\ & & qC^{i-1} \otimes_P \asymp & \xrightarrow{\quad} & qC^i \otimes_P \asymp & \xrightarrow{\quad} & qC^{i+1} \otimes_P \asymp \end{array} \quad (14)$$

¹Since the matching M will consist of only 0, 1, or 2 block isomorphisms per homological degree, one could avoid the graph theoretic language and simply perform Gaussian elimination on the morphisms of M one-by-one. The result will be identical, but our choice of notation will help us to analyze the result in Lemma 5.5.

where the maps S^i are diagonal matrices of saddles. In 14 the q :s at the bottom row denote the extra grading shifts associated to $\llbracket c \rrbracket_{\text{US}}$ and whereas the homological gradings and the internal gradings of C^i are hidden. Let us write the i :th chain space of C' as $D^i \oplus E^i$ where D^i is the direct sum of those diagrams which connect the two bottom strands and E^i is the direct sum of the rest:

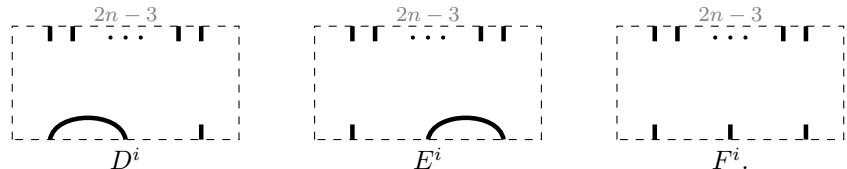


The diagrams of $D^i \otimes_P \asymp$ contain a circle, which can be delooped at the cost of duplicating them. More formally, there is an isomorphism $\begin{bmatrix} \mathfrak{d}_1 \\ \mathfrak{d}_2 \end{bmatrix} : qD^i \otimes_P \asymp \rightarrow q^2 \overline{D^{i-1} \otimes_P \asymp} \oplus \overline{D^{i-1} \otimes_P \asymp}$ where $\overline{D^i \otimes_P \asymp}$ denotes the direct sum of diagrams of $D^i \otimes_P \asymp$ with the new circle removed on every diagram, see Figure 2 for the local picture. Now decomposing C' and pulling back the differential we find an isomorphic complex C''

$$\begin{array}{ccccc}
 D^i \otimes_P \asymp & & D^{i+1} \otimes_P \asymp & & D^{i+2} \otimes_P \asymp \\
 \searrow & & \searrow & & \searrow \\
 E^i \otimes_P \asymp & \xrightarrow{m^i} & E^{i+1} \otimes_P \asymp & \xrightarrow{m^{i+1}} & E^{i+2} \otimes_P \asymp \\
 \dots & \xrightarrow{\overline{D^{i-1} \otimes_P \asymp}} & \xrightarrow{\overline{D^i \otimes_P \asymp}} & \xrightarrow{\overline{D^{i+1} \otimes_P \asymp}} & \dots \\
 q^2 \overline{D^{i-1} \otimes_P \asymp} & & q^2 \overline{D^i \otimes_P \asymp} & & q^2 \overline{D^{i+1} \otimes_P \asymp} \\
 qE^{i-1} \otimes_P \asymp & & qE^i \otimes_P \asymp & & qE^{i+1} \otimes_P \asymp
 \end{array}$$

where the chain space $(C'')^i$ is the direct sum of the 5 objects in the first column. The differentials of C'' are represented by 5×5 whose matrix element at $(1, 3)$ is m^i . Expanding the pullback differential gives $m^i = (-1)^i \mathfrak{d}_2(\pi_D \otimes_P \text{id})(\text{id} \otimes_P s)(\iota_D \otimes_P \text{id})$, where ι_D and π_D are the canonical inclusions and projections related to $C^i = D^i \oplus E^i$ and s is a saddle. Hence m^i are isomorphisms, $M = \{m^i \mid 0 \leq i \leq n(L)\}$ is a Morse matching and at Line 5 the Morse complex $M(C'') = C'''$ with $C''' \simeq C'$ is computed.

Type III: We decompose $C^i = D^i \oplus E^i \oplus F^i$ where D^i contains the diagrams which connect the two bottom left strands, E^i those which connect the two bottom right strands and F^i those which connects none of them which each other:



As with Type II, the summands of $D^i \otimes_P \asymp$ and $E^i \otimes_P \asymp$ contain circles which can be delooped so

that $C' = C \otimes_P \llbracket c \rrbracket_{\text{US}}$ is isomorphic to the complex C'' :

$$\begin{array}{ccc}
D^i \otimes_P) (& D^{i+1} \otimes_P) (& D^{i+2} \otimes_P) (\\
q^{-1} \overline{E^i \otimes_P} (& q^{-1} \overline{E^{i+1} \otimes_P} (& q^{-1} \overline{E^{i+2} \otimes_P} (\\
q \overline{E^i \otimes_P} (& q \overline{E^{i+1} \otimes_P} (& q \overline{E^{i+2} \otimes_P} (\\
F^i \otimes_P) (& F^{i+1} \otimes_P) (& F^{i+2} \otimes_P) (\\
\cdots & \cdots & \cdots \\
\overline{D^{i-1} \otimes_P \asymp} & \overline{D^i \otimes_P \asymp} & \overline{D^{i+1} \otimes_P \asymp} \\
m_1^i & m_2^i & m_2^{i+1} \\
q^2 \overline{D^{i-1} \otimes_P \asymp} & q^2 \overline{D^i \otimes_P \asymp} & q^2 \overline{D^{i+1} \otimes_P \asymp} \\
qE^{i-1} \otimes_P \asymp & qE^i \otimes_P \asymp & qE^{i+1} \otimes_P \asymp \\
qF^{i-1} \otimes_P \asymp & qF^i \otimes_P \asymp & qF^{i+1} \otimes_P \asymp
\end{array}$$

Again by expanding, one sees that the matrix elements m_1^i and m_2^i are isomorphisms. On the other hand, the dashed gray arrows are zero morphisms since the first tensor component of their expansions contains an inclusion $D^i \hookrightarrow D^i \oplus E^i$ followed by a projection $D^i \oplus E^i \twoheadrightarrow E^i$ (or alternatively $E^i \hookrightarrow D^i \oplus E^i \twoheadrightarrow D^i$). Hence $M = \{m_1^i, m_2^i \mid 0 \leq i \leq n(L)\}$ is a Morse matching, and at Line 5 we can perform Gaussian elimination on all of the morphisms m_1^i and m_2^i to obtain the Morse complex $M(C) = C'''$.

Type IV: Let us decompose $C^i = D^i \oplus E^i$ where D^i consists of those diagrams with \asymp connectivity and E^i of those with $) ($ connectivity. The compositions contain many new circles, but for the moment we only deloop the ones at the bottom of $D^i \otimes_P \asymp$ and the inner circle of $E^i \otimes_P) ($. This yields the complex C'' :

$$\begin{array}{ccc}
D^i \otimes_P) (& D^{i+1} \otimes_P) (& D^{i+2} \otimes_P) (\\
q^{-1} \overline{E^i \otimes_P} (& q^{-1} \overline{E^{i+1} \otimes_P} (& q^{-1} \overline{E^{i+2} \otimes_P} (\\
q \overline{E^i \otimes_P} (& q \overline{E^{i+1} \otimes_P} (& q \overline{E^{i+2} \otimes_P} (\\
F^i \otimes_P) (& F^{i+1} \otimes_P) (& F^{i+2} \otimes_P) (\\
\cdots & \cdots & \cdots \\
\overline{D^{i-1} \otimes_P \asymp} & \overline{D^i \otimes_P \asymp} & \overline{D^{i+1} \otimes_P \asymp} \\
m_1^i & m_2^i & m_2^{i+1} \\
q^2 \overline{D^{i-1} \otimes_P \asymp} & q^2 \overline{D^i \otimes_P \asymp} & q^2 \overline{D^{i+1} \otimes_P \asymp} \\
qE^{i-1} \otimes_P \asymp & qE^i \otimes_P \asymp & qE^{i+1} \otimes_P \asymp
\end{array}$$

from we set $M = \{m_1^i, m_2^i \mid 0 \leq i \leq n(L)\}$. The set M is a Morse matching since the gray dashed matrix elements on C''' are zero morphisms. To finish off with a simple complex, we deloop the rest

of the circles of $M(C)$ before calling the result C''' . (To more accurately follow the pseudocode of Algorithm 2 we could have delooped everything first before cancellation, but our equivalent two-step delooping process kept the notation lighter.)

5.3 Validity

The *truncation* of a chain complex (C, d) at $k \in \mathbb{Z}$, denoted by $\tau_{\leq k}(C)$, is the chain complex with differential ∂ defined as

$$\tau_{\leq k}(C)^i = \begin{cases} C^i, & \text{if } i \leq k \\ 0, & \text{otherwise} \end{cases} \quad \partial^i = \begin{cases} d^i, & \text{if } i < k \\ 0, & \text{otherwise.} \end{cases}$$

With this notation, Line 6 can be written as $C \leftarrow \tau_{\leq k+1}(C''')$. We say that chain complexes (C, d) and (D, ∂) over an additive category are *chain homotopic up to $k \in \mathbb{Z}$* , denoted $C \simeq_{\leq k} D$, if for all $i \leq k$ there are morphisms

$$f^i: C^i \rightarrow D^i, \quad g^i: D^i \rightarrow C^i, \quad s^i: C^i \rightarrow C^{i-1}, \quad t^i: D^i \rightarrow D^{i-1}$$

so that f^i and g^i commute with the differentials and s^i and t^i give the homotopy to the identities:

$$\begin{aligned} f^{i+1}d^i &= \partial^i f^i, & g^i f^i - \text{id}_{C^i} &= d^{i-1}s^i + s^{i+1}d^i \\ g^{i+1}\partial^i &= d^i g^i, & f^i g^i - \text{id}_{D^i} &= \partial^{i-1}t^i + t^{i+1}\partial^i. \end{aligned}$$

The usual definition of a chain homotopy is retrieved with $k = \infty$ and the following lemma is a collection of some useful basic properties.

Lemma 5.4. 1. $C \simeq D$ implies $C \simeq_{\leq k} D$ for all k .

2. $\tau_{\leq k+1}(C) \simeq_k C$, for all k .

3. If F is an additive functor, then $C \simeq_{\leq k} D$ implies $F(C) \simeq_{\leq k} F(D)$.

4. If C and D are chain complexes over an abelian category and $C \simeq_{\leq k} D$, then $H^i(C) \cong H^i(D)$ for all $i < k$.

5. Let P be a simple planar arc diagram, c a crossing, C and D complexes over $\text{Cob}_{\bullet/l}^{\mathbb{Z}}(2n)$ which are supported on the nonnegative homological degrees and $C \simeq_{\leq k} D$. Then $C \otimes_P \llbracket c \rrbracket_{\text{US}} \simeq_{\leq k} D \otimes_P \llbracket c \rrbracket_{\text{US}}$.

Proof. Claims 1, 2, and 3 follow directly from the definitions, 4 is a slight modification of the standard result and 5 uses the fact that $\llbracket c \rrbracket_{\text{US}}$ is supported in nonnegative homological degrees. \square

For analyzing correctness and estimating the time complexity of the algorithms \mathcal{A}_k we need to define the following *intermediate complexes*. We denote by $C_{[t]}$ the complex stored at variable C before running the for loop of \mathcal{A}_k (Lines 3-6 in Algorithm 2) for t :th time where $t \leq n(L) - 1$ and by $C_{[n(L)]}$ the complex which is stored at variable C after the for loop.

Proof that \mathcal{A}_k computes $\text{Kh}^{i,}(L)$ for $i \leq k - n_-(L)$ correctly.* At every line of Algorithm 2 one can use a corresponding claim of Lemma 5.4 to see that chain homotopy up to $k+1$ is preserved: At line 1 we initialize the complexes to be the same: $C_{[1]} = \llbracket c_1 \rrbracket_{\text{US}}$ so in particular $C_{[1]} \simeq_{\leq k+1} \llbracket T_1 \rrbracket_{\text{US}}$. Then to proceed with induction, we use Claim 5 at Line 3, Claim 1 at Lines 4 and 5, Claim 2 at Line 6 and Claim 3 at Line 8. This guarantees that

$$F(C_{[n]}) \simeq_{\leq k+1} F\llbracket L \rrbracket_{\text{US}} = u^{n_-(L)} q^{-n_+(L) + 2n_-(L)} \text{CKh}(L)$$

and hence by Claim 4 the algorithm \mathcal{A}_k computes $\text{Kh}^{i,*}(L)$ for $i \leq k - n_-(L)$ correctly. \square

5.4 Proof of polynomial time complexity for \mathcal{A}_k

In order to measure the sizes of objects and limit the number of operations, we establish two definitions. Firstly, for an object B in $\mathcal{Cob}_{\bullet/l}^{\mathbb{Z}}(2n)$, we set $\text{rank}_{\mathcal{Cob}}(B) = n$ where n is the number of loopless diagrams in the decomposition (10). Secondly, for a complex (C, d) over $\mathcal{Cob}_{\bullet/l}^{\mathbb{Z}}(2n)$, we define $\|C\|_i = \max\{\|f\| : f \text{ matrix element of } d^i\}$ for every integer i .

Lemma 5.5. *The following bounds hold for intermediate complexes $C_{[t]}$ of the algorithm \mathcal{A}_k*

$$\text{rank}_{\mathcal{Cob}}(C_{[t]}^i) \leq \binom{t}{i} \quad \text{for } t < n(L) \quad (15)$$

$$\text{rank}_{\mathcal{Cob}}(C_{[n(L)]}^i) \leq 2 \binom{n(L)}{i} \quad (16)$$

$$\|C_{[t]}\|_i \leq \exp_2 \left(\left(\frac{g}{2} + 1 \right) \left(\binom{t+1}{i+1} - 1 \right) \right) \quad \text{for all } t \quad (17)$$

where g is the girth of the nice scanning sequence and $\exp_2(a)$ denotes 2^a .

Proof. All three claims are proven with an induction on t . For (15) and (16) the base case $t = 1$ is trivial and the induction step splits into 4 simple cases depending on the type of the planar arc diagram P . For example, with Type II diagram and using the notation of Section 5.2 we can directly compute

$$\text{rank}_{\mathcal{Cob}}(C_{[t+1]}^i) = \text{rank}_{\mathcal{Cob}}(E^i \otimes_P) (\oplus q \overline{D^{i-1} \otimes_P} \widetilde{\circ} \oplus E^{i-1} \otimes_P \widetilde{\circ}) \leq \binom{t}{i} + \binom{t}{i-1} = \binom{t+1}{i}.$$

The additional factor 2 in (16) is associated to the extra deloopings involved with the Type IV diagram.

For (17) the base case $t = 1$ is routinely verified and we assume that (17) holds for some t . It is straightforward to see that

$$\|C_{[p]}\|_i = 0 \quad \text{when } i < 0 \text{ or } i \geq p \quad (18)$$

since the complex $\llbracket T_p \rrbracket$ is also 0 in homological degrees $i < 0$ and $i > p$. The morphisms of $C_{[t]}$ are also the morphisms of untruncated Morse complex $M(C'') = C'''$ which are sums of zig-zag paths. In the Morse complexes one can see that there are at most 3 paths on every sum and in Type III and IV they will use either $(m_1^i)^{-1}$, or $(m_2^i)^{-1}$ or neither but not both. The norm of each of the three summands will be bounded by $2^{g/2-1} \max(\|C_{[t]}\|_i, 1) \max(\|C_{[t]}\|_{i-1}, 1)$ where the $2^{g/2-1}$ term comes from Lemma 5.3. In particular we get

$$\|C_{[t+1]}\|_i \leq 3 \cdot 2^{\frac{g}{2}-1} \max(\|C_{[t]}\|_i, 1) \max(\|C_{[t]}\|_{i-1}, 1). \quad (19)$$

If $1 \leq i \leq t-1$, we can use the induction assumption and (19) to compute:

$$\begin{aligned} \|C_{[t+1]}\|_i &\leq 3 \cdot 2^{g-1} \exp_2 \left((g+1) \left(\binom{t+1}{i+1} - 1 \right) \right) \exp_2 \left((g+1) \left(\binom{t+1}{i} - 1 \right) \right) \\ &\leq \exp_2 \left((g+1) \left(\binom{t+2}{i+1} - 1 \right) \right). \end{aligned}$$

For $i = 0$ and $i = t$, the claim is obtained by using (18), (19), and the induction assumption. \square

Proof of the running time estimate for \mathcal{A}_k . As an input to \mathcal{A}_k , let us limit to link diagrams L with nice scanning sequences \mathcal{S} with a global upper bound $\text{girth}(\mathcal{S}) \leq g$ on them. By analyzing the pseudocode of Algorithm 2 line by line, we aim to show that on this fixed set of inputs \mathcal{A}_k runs in polynomial time with respect to $n(L)$.

By Lemma 5.5 the intermediate complex $C_{[t]}$ has at most $2 \sum_{j=0}^{k+1} \binom{t}{j} = \mathcal{O}(t^{k+1})$ objects. Since each matrix element of $C_{[t]}$ connects a pair of objects, there are at most $\mathcal{O}(t^{2k+2})$ matrix elements f and by Lemma 5.5 their norms are bounded by

$$\log_2 \|f\| \leq \max_{j \in \{0, \dots, k+1\}} \left(\binom{t+1}{j+1} - 1 \right) (g+1) = \mathcal{O}(t^{k+1}).$$

To generate $C' = C_{[t]} \otimes_{P_t} [c_{t+1}]_{\text{US}}$ from $C_{[t]}$ we perform at most $\mathcal{O}(t^{k+1})$ vertical compositions of objects, $\mathcal{O}(t^{2k+2})$ vertical compositions of morphisms and create at most $\mathcal{O}(t^{k+1})$ new saddle morphisms. When creating C'' from C' we deloop $\mathcal{O}(t^{k+1})$ objects and compose $\mathcal{O}(t^{2k+2})$ morphisms with the delooping $\mathfrak{d}_1, \mathfrak{d}_2$ and relooping $\mathfrak{r}_1, \mathfrak{r}_2$ morphisms. In constructing C''' from C'' we perform $\mathcal{O}(t^{2k+2})$ horizontal compositions and summations of morphisms. By Lemma 5.3 each horizontal composition takes polynomial time with respect to $\mathcal{O}(t^{k+1})$ since $\text{girth}(T_t) \leq \text{girth}(\mathcal{S}) \leq g$. The time taken for the truncation at Line 6 is negligible. We have now shown that there exists some constant $a \in \mathbb{Z}$, which is independent of t , and which ensures that only $\mathcal{O}(t^a)$ operations are performed at the t :th iteration of the for loop.

To execute the whole for loop, $n(L) - 1$ iterations are made with different $t \leq n(L) - 1$ values, so in particular at most

$$\sum_{t=1}^{n(L)-1} \mathcal{O}(t^a) \leq \sum_{t=1}^{n(L)-1} \mathcal{O}(n(L)^a) = \mathcal{O}(n(L)^{a+1})$$

operations are performed before the Smith normal form at Line 9. Smith normal form of an integer matrix $A^{n \times m}$ can be computed in polynomial time with respect to n, m and $\max_{i,j} \{\log_2(|A_{i,j}|)\}$, [Sto96]. By Lemma 5.5 the matrices of $F(C_{[n(L)]})$ admit polynomial upper bounds with respect to $n(L)$ for these quantities. Hence \mathcal{A}_k runs in polynomial time for the inputs (L, \mathcal{S}) with $\text{girth}(\mathcal{S}) \leq g$. Thus the running time of algorithm \mathcal{A}_k for all inputs (L, \mathcal{S}) is $\mathcal{O}(f(\text{girth}(\mathcal{S})) \cdot p_k(n(L)))$ for some function f and polynomial p_k . \square

6 Binomial rank bounds via nice scanning sequences

A graph G is *connected* if there is a path between every two vertices in G . In addition, a graph G is called *2-connected* if for every vertex $v \in G$ the induced graph of $G \setminus \{v\}$ is connected.

Lemma 6.1. *Let $G = (V, E)$ be a 2-connected graph and $A \subset V$ a subset of vertices so that $\#A \geq 1$ and $\#(V \setminus A) \geq 2$. If the induced graphs of A and $V \setminus A$ are connected there exists $x \in V \setminus A$ so that the induced subgraphs of $A \cup \{x\}$ and $V \setminus (A \cup \{x\})$ are also connected.*

Proof. Assume that x is a vertex of $V \setminus A$ for which the following maximum is obtained:

$$M = \max \{ \#C \mid x \in (V \setminus A), A \cup \{x\} \text{ connected, } C \text{ is a component of } V \setminus (A \cup \{x\}) \}.$$

If $M = \#V - \#A - 1$, then $V \setminus (A \cup \{x\})$ is connected so we can assume towards contradiction that $M < \#V - \#A - 1$.

Let C and D be two components of $V \setminus (A \cup \{x\})$ with $M = \#C$ and pick any vertices $c \in C$ and $d \in D$. Since G is 2-connected, there is a path P from d to c which does not cross x . Denote $d' \in D$ as the last vertex of P in D , so that $V \cup \{d'\}$ is connected and write C' for the component of c in $V \setminus (A \cup \{d'\})$. It follows that $\#C < \#C'$ since the component C' additionally contains the vertex x . On the other hand, $\#C' \leq M = \#C$ by maximality of M which gives us the desired contradiction. \square

A link diagram is called *reduced* if it does not contain any nugatory crossings, see Figure 7. It is straightforward to see that a connected link diagram L can be transformed into an isotopic, reduced, and connected link diagram L' with $n_+(L) \geq n_+(L')$ and $n_-(L) \geq n_-(L')$.



Figure 7: A nugatory crossing (left) and removing it (right). Removing the nugatory crossing flips the right hand side tangle T_2 by 180° around the x -axis which is indicated by the flipping the symbol T_2 to \mathfrak{T}^3 .

Lemma 6.2. *Every connected reduced link diagram admits a nice scanning sequence on it.*

Proof. Let L be a connected reduced link diagram and $G_L = (V, E)$ the 4-valent non-simple planar graph of L , whose vertices are the crossings of L and whose edges correspond to the strands between crossings. From the fact that G_L is connected and reduced, one can see that G_L is 2-connected. (Any vertex separator of G_L has to correspond to a nugatory crossing in L , but L does not have nugatory crossings since it is reduced.)

The sequence of crossings c_1, \dots, c_n for \mathcal{S} are retrieved from the graph G_L by a greedy algorithm. Start by picking any vertex $c_1 \in V$ and observe that the induced graphs $\{c_1\}$ and $V \setminus \{c_1\}$ are connected. Suppose next that c_1, \dots, c_i have been picked, so that $i \leq n-2$ and the induced subgraphs of $\{c_1, \dots, c_i\}$ and $V \setminus \{c_1, \dots, c_i\}$ are connected. By Lemma 6.1 there is a vertex c_{i+1} for which that the induced subgraphs of $\{c_1, \dots, c_{i+1}\}$ and $V \setminus \{c_1, \dots, c_{i+1}\}$ also remain connected. Finally, once c_1, \dots, c_{n-1} have been chosen, add the last vertex in to complete the sequence of crossings c_1, \dots, c_n .

The sequence of subtangle diagrams T_1, \dots, T_n and simple planar diagrams P_1, \dots, P_{n-1} can be retrieved by setting $T_1 = c_1$ and analyzing how c_{i+1} inductively glues into T_i . The induced graph of c_1, \dots, c_i remains connected after gluing in c_{i+1} which is why P_i has to connect c_{i+1} to T_i with 1, 2, 3, or 4 strands. Moreover the vertex c_{i+1} cannot be connected to T_i with 4 strands except when $i+1 = n$, since the induced graph of $\{c_{i+1}, \dots, c_n\}$ is connected. Finally, no strand can go from a crossing to itself since L was reduced. Thus c_{i+1} is glued to T_i with a Type I, II, or III simple planar arc diagram, when $i \leq n-2$ and with a Type IV diagram when $i = n-1$ and so $\mathcal{S} = ((T_1, \dots, T_n), (c_1, \dots, c_n), (P_1, \dots, P_{n-1}))$ forms a nice scanning sequence for L . \square

Proof of Proposition 1.5. Let L' denote the link diagram which is obtained from L by removing the nugatory crossings. By Lemma 6.2 there exists a nice scanning sequence \mathcal{S} on L' . When running the algorithm A_k on (L', \mathcal{S}) with $k = n(L')$ no nonzero chain spaces get truncated so there is an honest chain homotopy $u^{-n_-(L)} q^{n_+(L)-2n_-(L)} C_{[n(L)]} \simeq \llbracket L' \rrbracket$. Hence for any field \mathbb{F} we can compute

$$\begin{aligned} \dim_{\mathbb{F}}(\text{Kh}^{i,*}(L; \mathbb{F})) &= \dim_{\mathbb{F}}(\text{Kh}^{i,*}(L'; \mathbb{F})) \leq \text{rank} \left(F(u^{-n_-(L)} q^{n_+(L)-2n_-(L)} C_{[n(L)]})^i \right) \\ &\leq \text{rank}_{\mathcal{C}ob} \left(C_{[n(L)]}^{i+n_-(L)} \right) \leq 2 \binom{n(L')}{i+n_-(L')} \leq 2 \binom{n(L)}{i+n_-(L)} \end{aligned}$$

by bounding the homology with the size of the chain complex and using Lemma 5.5. \square

7 Asymptotical strictness of rank bounds

Proposition 1.6 is a non-vanishing result which claims an asymptotic lower bound on the the rank of certain Khovanov homology groups. To prove it we consider the chain complexes $\text{CKh}(L_t)$, where L_t is the braid closure of the braid diagram $(\sigma_1 \sigma_3 \sigma_2^4)^t \sigma_1 \sigma_3$, see Figure 8. On $\text{CKh}(L_t)$ we employ an algorithmic Morse matching M_{lex} from the previous work of the first author [Kel25]. This produces a Morse complex $M_{\text{lex}}(\text{CKh}(L_t))$ where we find 2^t explicit, distinct, non-vanishing homology cycles.

Since we are using M_{lex} on link diagrams L_t which only contain crossings of the form \times the construction of M_{lex} can be simplified quite a bit. In [Kel25] the matching M_{lex} was defined on Bar-Natan's tangle complexes but since we only work with specific links, we can define it directly for $\text{CKh}(L_t)$ which are chain complexes of free \mathbb{Z} -modules. The direct sum decomposition which we fix on $\text{CKh}(L_t)$ consists of single copies of \mathbb{Z} : one for each 2-coloring of each smoothing of L_t . The two grayscale colors we use are dashed gray \circlearrowleft for the unit 1 and thick black \bullet for the counit x . This allows us to write merge m and split Δ maps of the Khovanov complex as

$$\begin{aligned} m(\circlearrowleft \otimes \circlearrowleft) &= \circlearrowleft & \Delta(\circlearrowleft) &= \circlearrowleft \otimes \bullet + \bullet \otimes \circlearrowleft \\ m(\bullet \otimes \bullet) &= m(\circlearrowleft \otimes \bullet) = \bullet & \Delta(\bullet) &= \bullet \otimes \bullet \\ m(\bullet \otimes \bullet) &= 0. \end{aligned}$$

In particular, splitting out \bullet will always be an isomorphism and our Morse matching M_{lex} will consist of such splits.

We order the crossings of L_t from bottom to top or equivalently the characters of $(\sigma_1 \sigma_3 \sigma_2^4)^t \sigma_1 \sigma_3$ from left to right. For a matrix element f of $\text{CKh}(L_t)$ which changes a 0 smoothing to a 1 smoothing at the k :th crossing, we denote $i(f) = k$. Diagrammatically, we define

$$M = \left\{ f \in \{\text{matrix elements of } \text{CKh}(L_t)\} \mid \begin{array}{l} f \text{ splits out a } \bullet \text{ loop which completely located} \\ \text{at or below the } i(f)\text{:th crossing in the diagram } L_t \end{array} \right\}.$$

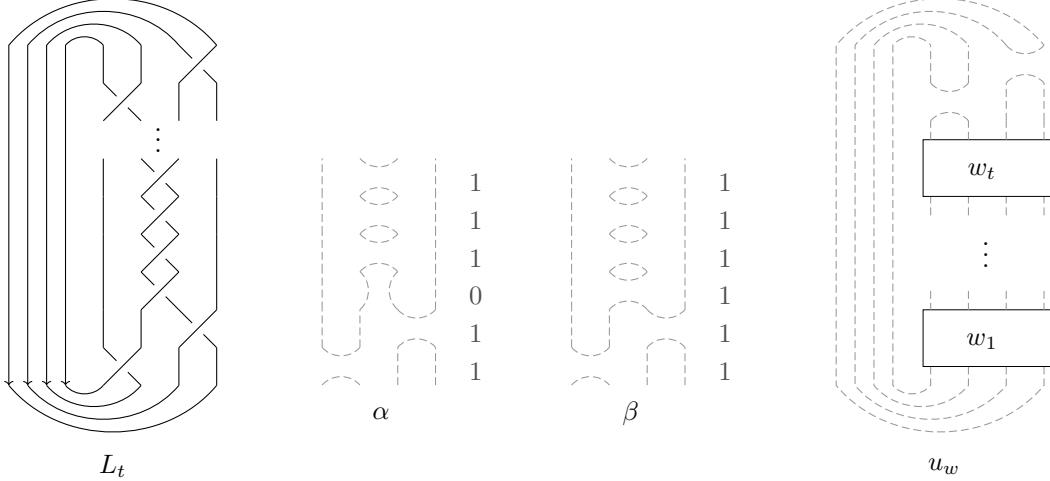


Figure 8: The link diagram L_t and local pictures α, β . The special vertices u_w of $G(\text{CKh}(L_t), M_{\text{lex}})$ are constructed by taking a tuple $w \in \{\alpha, \beta\}^t$ and gluing the local pictures of w on top of each other, 1 smoothing the two last vertices and using the dashed gray color on all loops.

For every vertex x we define $N(x)$ to be the set of neighboring morphisms from M :

$$N(x) = \{f: a \rightarrow x \mid f \in M\} \cup \{g: x \rightarrow b \mid g \in M\}.$$

This allows us to finally set

$$M_{\text{lex}} = \{f: x \rightarrow y \mid f \in M, i(f) \leq i(g) \text{ for all } g \in N(x) \cup N(y)\}.$$

Proposition 3.5 of [Kel25] shows that M_{lex} is a Morse matching for every tangle/link diagram and in particular:

Lemma 7.1. *The matching M_{lex} is a Morse matching on $\text{CKh}(L_t)$.*

In Figure 8 we construct special vertices u_w of $G(\text{CKh}(L_t), M_{\text{lex}})$ for every $w \in \{\alpha, \beta\}^t$. These pictures u_w end up representing distinct homology cycles in the Morse complex $M_{\text{lex}}(\text{CKh}(L_t))$. The proof of Proposition 1.6 comes as a consequence of the following three claims on the diagrams u_w :

1. The vertex u_w is unmatched in M_{lex} and thus it represents a copy of \mathbb{Z} in homological degree $6t + 2 - \#\{j \mid w_j = \alpha\}$ of the Morse complex $M_{\text{lex}}(\text{CKh}(L_t))$.
2. Suppose the homological degree of u_w is h and denote the projection of $(M_{\text{lex}}(L_t))^h$ to u_w by π . Then $\pi\partial^{h-1} = 0$.
3. Suppose the homological degree of u_w is h and denote the inclusion of u_w into $(M_{\text{lex}}(L_t))^h$ by ι . Then $\partial^h \iota = 0$.

Proof of Proposition 1.6. By Claim 1 every u_w represents a copy of \mathbb{Z} in the complex $M_{\text{lex}}(\text{CKh}(L_t))$ and by Claims 2 and 3 that \mathbb{Z} fully contributes to the homology $H(M_{\text{lex}}(\text{CKh}(L_t))) \cong \text{Kh}(L_t)$. This yields

$$\text{Kh}^{n(L_t)-k,*}(L_t) \cong \mathbb{Z}^{\binom{t}{k}} \oplus A_{k,t}$$

for some \mathbb{Z} -modules $A_{k,t}$. Thus

$$\liminf_{t \rightarrow \infty} \left(\frac{\text{rank}(\text{Kh}^{n(L_t)-k,*}(L_t))}{\binom{n(L_t)}{n(L_t)-k}} \right) \geq \liminf_{t \rightarrow \infty} \left(\frac{\binom{t}{k}}{\binom{6t+2}{6t+2-k}} \right) = \frac{1}{6^k}$$

concluding the proof. \square

Proof of Claim 1. Changing any 0-smoothing of u_w to a 1-smoothing generates a split morphism $f: u_w \rightarrow y$. No matter what crossing we chose to act on by f , neither of the loops in y which are split by f will be completely contained below $i(f)$ in y . Hence $f \notin M$. On the other hand, changing any 1-smoothing of u_w to a 0-smoothing is geometrically a split in the reverse direction. However u_w does not contain any \bullet loops and every split morphisms generates at least one \bullet loop, so there cannot exist an edge $g: x \rightarrow u_w$ in the graph $G(\text{CKh}(L_t))$. We conclude that $N(u_w) = \emptyset$ and so u_w is unmatched. \square

Proof of Claim 2. Let x be an unmatched vertex of $G(\text{CKh}(L_t), M_{\text{lex}})$ with homological degree $h-1$. In the previous claim we argued that there are no edges in $x \rightarrow u_w$ in the graph $G(\text{CKh}(L_t))$. It follows that cannot exist any zig-zag paths from x to u_w . Thus $\partial_{u_w, x} = 0$ and so $\pi\partial^{h-1} = 0$ as well. \square

7.1 Proof of Claim 3

Let y be an unmatched vertex of $G(\text{CKh}(L_t), M_{\text{lex}})$ with homological degree $h+1$ and set $P = \{\text{paths: } u_w \rightarrow y\}$. In Claim 2 we deduced $\pi\partial^{h-1} = 0$ from the fact that there were no paths nontrivial paths from any vertex to u_w in $G(\text{CKh}(L_t), M_{\text{lex}})$. Since P can be non-empty the same tactic cannot be used to prove $\partial^h = 0$. Instead, we will show that P can be arranged into distinct pairs $(p, \lambda(p))$ for which $R(\lambda(p)) = -R(p)$. From this pairing, one can easily compute:

$$\partial_{y, u_w} = \sum_{q \in P} R(q) = \sum_{(p, \lambda(p))} R(p) + R(\lambda(p)) = 0.$$

We say that $l_1 \rightarrow \dots \rightarrow l_8$ is an l -subpath if when cutting the diagrams of l_1, \dots, l_8 above the $6n+8$:th crossing for some $0 \leq n \leq t-1$ generates the pictures in Figure 9c. Similarly, we call $s_1 \rightarrow s_2$ a s -subpath if a similar cut generates diagrams of Figure 9b. In order to prove that every $p \in P$ contains either an s - or an l -subpath we will need the following result.

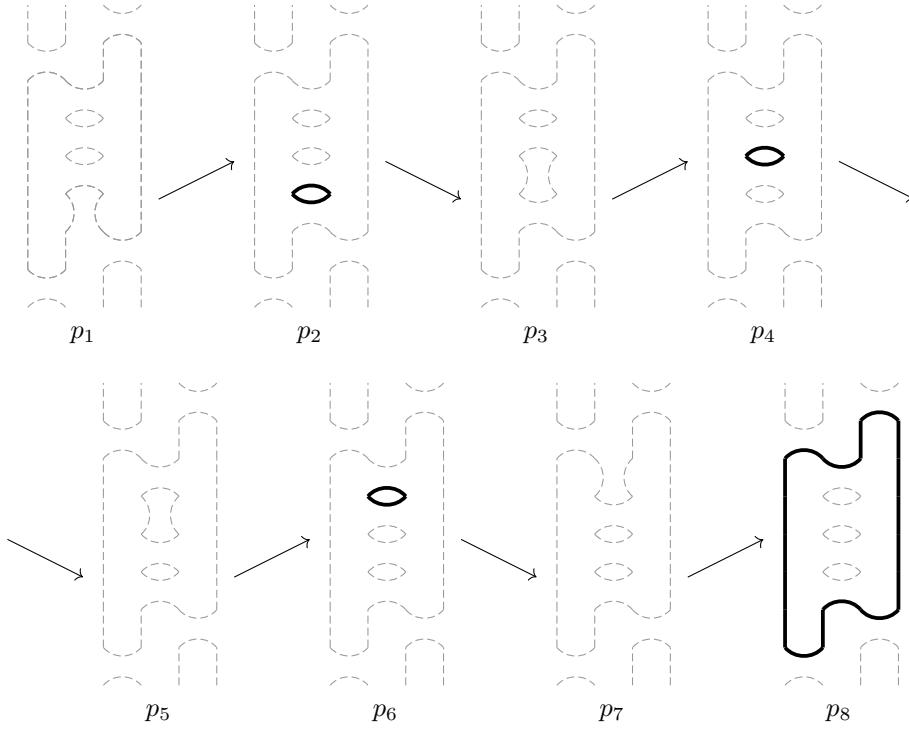
Lemma 7.2 (Lemma 3.4 from [Kel25]). *Let $(b \rightarrow a) \in M_{\text{lex}}$ and $a \rightarrow b \rightarrow c$ be a subpath of a zig-zag path between unmatched cells in the graph $G(\text{CKh}(L_t), M_{\text{lex}})$. Then $i(b \rightarrow a) \leq i(b \rightarrow c)$.*

Let $p = (v_1 \rightarrow \dots \rightarrow v_{2m}) \in P$ be a zig-zag path which does not contain an s -subpath. By observing the diagram of u_w one can see that the $v_1 \rightarrow v_2$ has to be of the form $s_1 \rightarrow s_2$ from Figure 9b or $l_1 \rightarrow l_2$ from Figure 9c. Since p does not have an s -subpath the latter must hold. Next, let $v_{2k+1} \rightarrow v_{2k+2}$ be an edge in p which is of the form $l_1 \rightarrow l_2$ and for which $i(v_{2k+1} \rightarrow v_{2k+2}) \leq i(v_{2j+1} \rightarrow v_{2j+2})$ for every other edge $v_{2j+1} \rightarrow v_{2j+2}$ in p of the form $l_1 \rightarrow l_2$. By verifying that vertices of type l_2 are matched to vertices of type l_3 we obtain that v_{2k+3} is of type l_3 . From Lemma 7.2 we get $i(v_{2k+3} \rightarrow v_{2k+4}) \leq i(v_{2k+3} \rightarrow v_{2k+2})$. The only edges from v_{2k+3} with i value strictly below $i(v_{2k+3} \rightarrow v_{2k+2})$ are either of the type $s_1 \rightarrow s_2$ or $l_1 \rightarrow l_2$ but neither of these are possible due to p not containing an s -subpath and the minimality of $i(v_{2k+1} \rightarrow v_{2k+2})$. It follows that $i(v_{2k+3} \rightarrow v_{2k+4}) = i(v_{2k+3} \rightarrow v_{2k+2})$ and so v_{2k+4} must be of type l_4 (as the arrow to type l_2 vertex is reversed). By alternately repeating the two aforementioned arguments one can deduce that $v_{2k+1} \rightarrow \dots \rightarrow v_{2k+8}$ is an l -subpath. Hence every $p \in P$ contains an s - or an l -subpath.

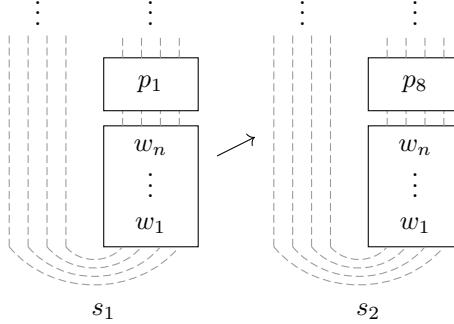
The pairing we foreshadowed takes the form of a bijection $\lambda: P \rightarrow P$. The function λ takes in a path p , cuts out the first occurrence of an s - or l -subpath and glues back in an l - or an s -subpath. If s -subpath was cut out, then l -subpath is glued in and vice versa. Since s - and l -subpaths begin and end at similar vertices, the map λ is a well defined and since $\lambda(\lambda(p)) = p$, it is a bijection.

To confirm that $R(\lambda(p)) = -R(p)$ we take a path $p \in P$ where an s -path occurs first. Denote $s_1 \rightarrow s_2$ and $l_1 \rightarrow \dots \rightarrow l_8$ as the cut out s -path of p and the l -path which is glued back in by λ . Perhaps surprisingly, we do not need to fix a sign convention on the complex $\text{CKh}(L_t)$ to see that the signs of $R(s_1 \rightarrow s_2)$ and $R(l_1 \rightarrow l_2)$ agree with each other. On the other hand, the signs of $R(l_2 \rightarrow l_3)$ and $R(l_3 \rightarrow l_4)$ disagree with each other as do those of $R(l_4 \rightarrow l_5)$ and $R(l_5 \rightarrow l_6)$ and further those of $R(l_6 \rightarrow l_7)$ and $R(l_7 \rightarrow l_8)$. Putting all of this together gives $R(s_1 \rightarrow s_2) = -R(l_1 \rightarrow \dots \rightarrow l_8)$ and since the paths agree elsewhere we also get $R(\lambda(p)) = -R(p)$ which concludes the proof. \square

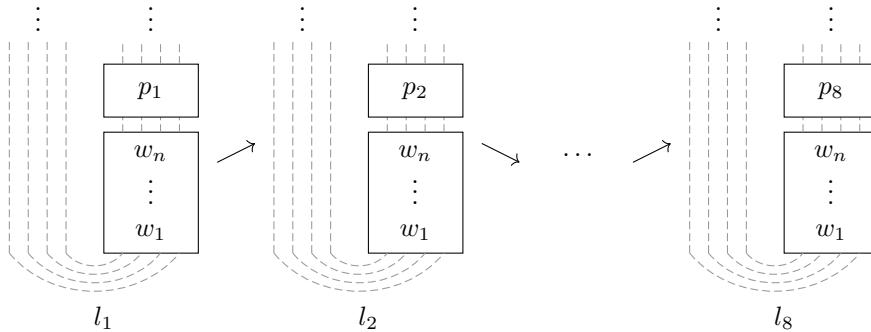
Remark 7.3. In proving Proposition 1.6, the signs were only needed at the last step. There, it sufficed that every edge $e: a \rightarrow b$ of the hypercube $\text{CKh}(L_t)$ between two smoothings a and b has a consistent sign no matter what colors are chosen on the loops of a and b . This fact also holds for all sign patterns of odd Khovanov homology [ORS13]. Hence the proof can be carried out in the odd setting and Proposition 1.6 holds for odd Khovanov homology as well.



(a) The local pictures p_1, \dots, p_8 which are used to make the s - and l -subpaths.



(b) An s -subpath. The lower boxes $(w_1, \dots, w_n) \in \{\alpha, \beta\}^n$ contain n -copies of the α and β pictures from Figure 8 stacked on top of each other.



(c) An l -subpath. The lower boxes $(w_1, \dots, w_n) \in \{\alpha, \beta\}^n$ contain n -copies of the α and β pictures from Figure 8 stacked on top of each other.

Figure 9: The s -subpaths, Figure 9b, and l -subpaths, Figure 9c, are paths in $G(\text{CKh}(L_t), M_{\text{lex}})$. The vertices of s - and l -subpaths are global diagrams whose bottom halves match the pictures of 9b and 9c.

Remark 7.4. While much of the construction of the non-vanishing homology cycles C_w was local, the global requirement that all neighboring morphisms of C_w were splits was also essential. This leaves some room for tweaking the construction of L_t . By considering braid closures L'_t of $(\sigma_1\sigma_3\sigma_2^4)^{2t}\sigma_1\sigma_3\sigma_2^3$ one can build similar homology cycles C'_w by 1-smoothing the uppermost 5 crossings. It follows that the asymptotic lower bounds of Proposition 1.6 can be realized with knots and therefore it is not a special feature of 2 or 4 component links.

Typically, the Jones polynomial is much easier to obtain than Khovanov homology both theoretically and computationally. In the proof of Theorem 1.3 we already used this fact to our advantage when obtaining the Khovanov homology of quasi-alternating links from the Jones polynomial and signature. In Proposition 1.6 the situation is remarkably opposite: we obtain lower bounds on the ranks of Khovanov homology groups, which lie on a single diagonal, without being able to say anything about the Jones polynomials $J_{L_t}(q)$. Since the links L_t are not homologically thin for t large enough², we cannot ensure that in decatgeorifying $\text{Kh}(L_t)$ to $J_{L_t}(q)$ the monomials corresponding to the high rank homology groups will not get canceled in the alternating sum.

8 Open questions

In [PS24], Przytycki and Silvero show that the extremal Khovanov homology of a closed 4-braid L can be computed in polynomial time. Furthermore, there are at most two non-trivial groups $\text{Kh}^{i, q_{\min}}(L)$ and $\text{Kh}^{k, q_{\min}}(L)$ with $k \geq i$, both free abelian and total rank at most 4 [PS24, Cor.1.6]. A priori, the gap $k - i$ could be arbitrarily large, so our methods do not give a direct proof of their result. Still, one can ask whether a variation of our Algorithm 2 which only keeps track of minimal quantum degrees can be used to recover their result.

Question 8.1. Is there a modification of Algorithm 2 which calculates the extremal Khovanov homology of a closed k -braid in polynomial time?

The scanning algorithm of Bar-Natan glues complexes of single crossings with simple planar arc diagrams of Type I-IV before simplifying the total complex with iterated Gaussian elimination. If implemented naively, this procedure requires one to build a larger complex before starting the simplification process:

$$C \xrightarrow{\text{glue}} C \otimes_P \llbracket c \rrbracket \xrightarrow{\text{deloop}} D \xrightarrow{\substack{\text{iterate Gaussian} \\ \text{elimination}}} E.$$

The Morse complexes we build for Type II, III, and IV diagrams and crossings $c \in \{\times, \times\}$ in Section 5.2 could be written explicitly and implemented matrix element by matrix element without first delooping, which accounts to

$$C \xrightarrow{\substack{\text{build the Morse complex} \\ \text{associated to } c \text{ and } P}} M(C) \xrightarrow{\text{iterate Gaussian elimination}} E.$$

It is straight-forward to see that this alternative, more laborious, workflow will still take exponential time on knots with exponentially large Betti numbers. Nevertheless, there ought to be room for optimization and the following question could be investigated empirically.

Question 8.2. Would implementing a Morse theoretic simplification as part of the scanning algorithm significantly speed up the computation for average-case knots?

In addition to average-case knots, one could hope that computing Khovanov homology would be fast for links of special interest. In the study of Khovanov homology, torus links have inarguably received particular attention and it is natural to ask:

Question 8.3. Let $n \geq 4$. Does Bar-Natan's scanning algorithm compute the Khovanov homology of an n -strand torus link in polynomial time, either with integral or field coefficients?

Bar-Natan's scanning algorithm computes the Khovanov homology of 2- and 3-strand torus links in polynomial time. This can be seen from the techniques in [Sch25b, §3,4]. For $n = 4$, the Bar-Natan complexes $\llbracket (\sigma_1\sigma_2\sigma_3)^n \rrbracket$ of open braids $(\sigma_1\sigma_2\sigma_3)^n$ can be simplified to complexes with quadratically many generators [Kel25, Corollary 4.2] which supports a positive answer to Question 8.3, at least

²By a result of Islambouli and Willis [IW18], the complexes of positive braids $(\sigma_1\sigma_3\sigma_2^4)^t\sigma_1\sigma_3$ converge to the categorified Jones-Wenzel projector. Hence the Khovanov homology of links L_t , suitably normalized, converges to the stable Khovanov homology of $T(4, \infty)$. Since the $\text{Kh}(T(4, \infty))$ is not homologically thin in the small (≤ 4) homological degrees ([Sto07] Theorem 3.4) neither can L_t be for large enough t .

when using coefficients from a finite field. An affirmative answer is also suggested by the following proposition, the proof of which was kindly provided to the authors by Qiuyu Ren.

Proposition 8.4. *There exist $\{c_n\}_{n \geq 2}$, such that $\dim_{\mathbb{F}}(\text{Kh}(T(n, m); \mathbb{F})) \leq c_n m^{\lfloor \frac{n}{2} \rfloor}$ for all fields \mathbb{F} .*

Proof. Let $L_{n,k}$ denote the braid closure of $(\sigma_1 \dots \sigma_{n-1})^m \sigma_1 \dots \sigma_r$ where $k = (n-1)m + r$ and $0 \leq r < n-1$. Since $T(n, m) = L_{n, (n-1)m}$, it suffices to show that there exists coefficients c_n , such that $\dim_{\mathbb{F}}(\text{Kh}(L_{n,m}; \mathbb{F})) \leq c_n m^{\lfloor \frac{n}{2} \rfloor}$. We proceed with an induction on n ; sufficient coefficients c_2 and c_3 can be obtained from [Kho00, §6.2], [CLSS22, Cor.5.7] and Theorem 3.3. Assume next that there exists c_2, \dots, c_{n-1} with $\dim_{\mathbb{F}}(\text{Kh}(L_{i,m}; \mathbb{F})) \leq c_i m^{\lfloor \frac{i}{2} \rfloor}$.

The last crossing of $L_{n,m}$ can be resolved in two ways, which generates two links $L_{n,m}[\square]$ and $L_{n,m}[\bigcirc]$. The link $L_{n,m}[\square]$ is isotopic to $L_{n,m-1}$ whereas $L_{n,m}[\bigcirc]$ is isotopic to some $L_{s,t}$ or $L_{s,t} \sqcup$ (unknot) for some $s \leq n-2$ and $t \leq m$. The isotopy argument for $L_{n,m}[\bigcirc]$ is explained well in [Tag17, Figures 15-20] although our conventions are mirrored.

Associated to the last crossing of the braid word $L_{n,m}$, there is an exact triangle

$$\begin{array}{ccc} \text{Kh}(L_{n,m}[\bigcirc]; \mathbb{F}) & \longrightarrow & \text{Kh}(L_{n,m}; \mathbb{F}) \\ & \swarrow & \downarrow \\ & \text{Kh}(L_{n,m}[\square]; \mathbb{F}) & \end{array}$$

When plugging in the aforementioned isotopies, this yields an inequality on the dimensions:

$$\dim_{\mathbb{F}}(\text{Kh}(L_{n,m}; \mathbb{F})) \leq \dim_{\mathbb{F}}(\text{Kh}(L_{n,m-1}; \mathbb{F})) + 2 \dim_{\mathbb{F}}(\text{Kh}(L_{s,t}; \mathbb{F})).$$

Hence choosing $c_n = 2 \max_{i < n} c_i$ gives the desired bound for n . \square

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