# The worst data for hierarchical log-linear models 

František Matúš (Prague)

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## The hierarchical model $\mathcal{E}_{N, A}$

Let $N$ be a nonempty finite set,
$\mathcal{A}$ a family of subsets of $N$ such that $\bigcup \mathcal{A}=N$, and $X=\prod_{i \in N} X_{i}$ the Cartesian product of finite state spaces.

A probability measure (pm) $Q$ on $X$ is called $\mathcal{A}$-factorizable if for each $I \in \mathcal{A}$ there exists a real function $\psi_{I}$ on $X_{I}=\prod_{i \in I} X_{i}$ s.t.

$$
Q(x)=\prod_{l \in \mathcal{A}} \psi_{l}\left(\pi_{l} x\right), \quad x \in X
$$

where $\pi_{l}$ projects $x$ to $X_{l}$.
The set of all $\mathcal{A}$-factorizable pm's that are positive, $Q(x)>0$ for $x \in X$, is denoted by $\mathcal{E}_{N, \mathcal{A}}$.

## Information divergence from a model

The information divergence or relative entropy between pm's $P, Q$ on $X$ is given by

$$
D(P \| Q)= \begin{cases}\sum_{x: P(x)>0} P(x) \ln \frac{P(x)}{Q(x)}, & \text { if } P \ll Q \\ +\infty, & \text { otherwise }\end{cases}
$$

and the divergence of $P$ from a model $\mathcal{E}$ by

$$
D(P \| \mathcal{E})=\inf _{Q \in \mathcal{E}} D(P \| Q)
$$

If $P$ is the empirical distribution of a dataset then a miminizer $Q$ corresponds to an MLE estimate from the data in the model $\mathcal{E}$.

The number $D(P \| \mathcal{E})$ characterizes fit of the data to the model.

## The worst data

The problem of maximization

$$
\max \{D(P \| \mathcal{E}): P \text { pm on } X\}
$$

goes back to Nihat Ay (2004) Ann. Probab.
A maximizer $P$ admits interpretation as the empirical distribution of a bad dataset.

Example:
$\mathcal{E}=\operatorname{Bi}(n), n \geqslant 3$, has the unique global maximizer $\frac{1}{2}\left(\delta_{0}+\delta_{n}\right)$.
In general difficult, even for 4 binary variables with all 2-way interactions.

## Upper bound on the divergence from $\mathcal{E}_{N, A}$

## Theorem

For any pm P on $X$

$$
D\left(P \| \mathcal{E}_{N, \mathcal{A}}\right) \leqslant \min _{I \in \mathcal{A}} \sum_{i \in N \backslash I} H\left(\pi_{i} P\right)
$$

(Shannon entropies of marginals)
Proof: induction on $|N|$, decomposition tricks, ...
As a consequence, assuming all spaces $X_{i}$ of the cardinality $d$,

$$
\max D\left(\cdot \| \mathcal{E}_{N, \mathcal{A}}\right) \leqslant \min _{l \in \mathcal{A}} \sum_{i \in N \backslash I} \ln \left|X_{i}\right| \leqslant\left[|N|-\max _{I \in \mathcal{A}}|I|\right] \ln d
$$

For 4 binary variables with all 2-way interactions the bound $2 \ln 2$ is, however, not tight.

## Matroidal hierarchical models

Consider a simple connected matroid with the ground set $N$ of the cardinality $n$, the rank function $r$, and the family of bases $\mathcal{A} \subseteq\binom{N}{k}$. Let all state spaces $X_{i}$ have the same cardinality $d$.

## Theorem

If a pm $P$ on $X$ satisfies

$$
H\left(\pi_{l} P\right)=r(I) \ln d, \quad I \subseteq N
$$

then it attains the upper bound, $D\left(P \| \mathcal{E}_{N, \mathcal{A}}\right)=[n-k] \ln d$.
The converse holds if the matroid is uniform.
This set of equalities is equivalent to saying that $P$ is an ideal secret sharing scheme (sss) on the set of participants $N$ with any choice of the dealer $i \in N$ and a secret of size $d$
(an object studied in cryptography for more than two decades).

Example: all $k$-way interactions, $\mathcal{A}=\binom{N}{k}$, among $n$ variables, each taking $d$ values. An ideal sss corresponds to an ( $n-k$ )-tuple of orthogonal Latin hypercubes of the size $d$.

## CONCLUSION

Data remote to a statistical model can have a distinct cryptographic meaning.

